

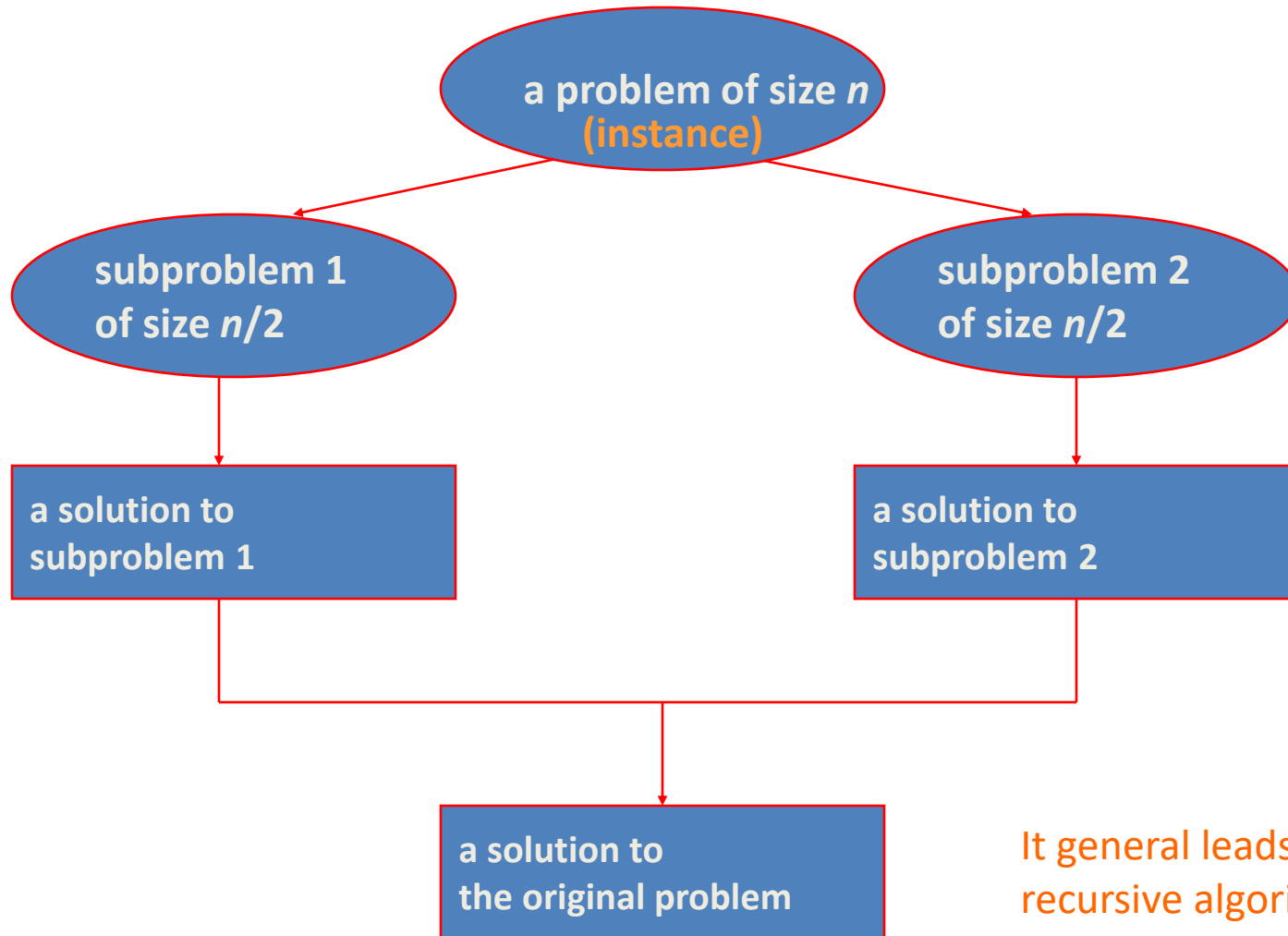
Divide and Conquer Algorithms

Divide-and-Conquer

The most-well known algorithm design strategy:

1. Divide instance of problem into two or more smaller instances
2. Solve smaller instances recursively
3. Obtain solution to original (larger) instance by combining these solutions

Divide-and-Conquer Technique (cont.)



It general leads to a recursive algorithm!

Linear Search

Problem: Given a list of N values, determine whether a given value X occurs in the list.

For example, consider the problem of determining whether the value 55 occurs in:

1	2	3	4	5	6	7	8
17	31	9	73	55	12	19	7

There is an obvious, correct algorithm:

start at one end of the list,
if the current element doesn't equal the search target, move to the next element,
stopping when a match is found or the opposite end of the list is reached.

Basic principle: divide the list into the current element and everything before (or after) it; if current isn't a match, search the other case

Linear Search

```
algorithm LinearSearch takes number X, list number L, number Sz

# Determines whether the value X occurs within the list L.
# Pre: L must be initialized to hold exactly Sz values
#

# Walk from the upper end of the list toward the lower end,
# looking for a match:

while Sz > 0 AND L[Sz] != X
    Sz := Sz - 1
endwhile

if Sz > 0          # See if we walked off the front of the list
    display true   # if not, got a match
else
    display false # if so, no match

halt
```

Binary Search

Repeatedly halving the size of the “search space” is the main idea behind the method of **binary search**.

An item in a sorted array of length **n** can be located with just **$\log_2 n$** comparisons.

Binary Search

Repeatedly halving the size of the “search space” is the main idea behind the method of **binary search**.

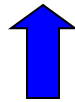
An item in a sorted array of length **n** can be located with just **$\log_2 n$** comparisons.

“Savings” is significant!

n	$\log_2(n)$
100	7
1000	10
10000	13

Binary search: target $x = 70$

	1	2	3	4	5	6	7	8	9	10	11	12
v	12	15	33	35	42	45	51	62	73	75	86	98



L: 1

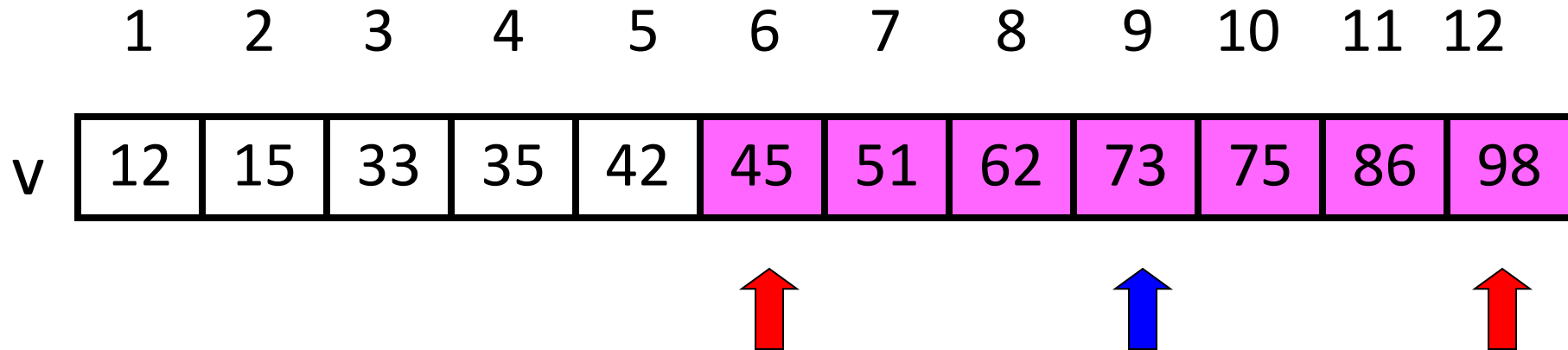
$v(\text{Mid}) \leq x$

Mid: 6

R: 12

So throw away the left half...

Binary search: target $x = 70$



L: 6

Mid: 9

R: 12

$x < v(\text{Mid})$

So throw away the right half...

Binary search: target $x = 70$

	1	2	3	4	5	6	7	8	9	10	11	12
v	12	15	33	35	42	45	51	62	73	75	86	98



L:

6

$v(\text{Mid}) \leq x$

Mid:

7

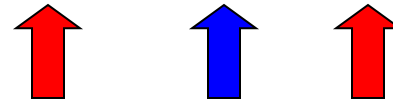
So throw away the left
half...

R:

9

Binary search: target $x = 70$

	1	2	3	4	5	6	7	8	9	10	11	12
v	12	15	33	35	42	45	51	62	73	75	86	98



L:

7

$v(\text{Mid}) \leq x$

Mid:

8

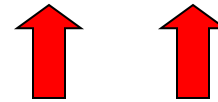
So throw away the left half...

R:

9

Binary search: target $x = 70$

	1	2	3	4	5	6	7	8	9	10	11	12
v	12	15	33	35	42	45	51	62	73	75	86	98



L:

8

Mid:

8

R:

9

Done because

$$R - L = 1$$

```

function L = BinarySearch(a,x)
# x is a row n-vector with x(1) < ... < x(n)
# where x(1) <= a <= x(n)
# L is the index such that x(L) <= a <= x(L+1)

L = 1;  R = length(x);
# x(L) <= a <= x(R)
while R-L > 1
    mid = floor((L+R)/2);
    # Note that mid does not equal L or R.
    if a < x(mid)
        # x(L) <= a <= x(mid)
        R = mid;
    else
        # x(mid) <= a <= x(R)
        L = mid;
    end
end
end

```

Binary search is efficient, but how do we sort a vector in the first place so that we can use binary search?

- Many different algorithms out there...
- Let's look at **merge sort**
- An example of the “divide and conquer” approach

Merge sort: Motivation

If I have two helpers, I'd...

- Give each helper half the array to sort
- Then I get back the sorted subarrays and **merge** them.

What if those two helpers each had two sub-helpers?

And the sub-helpers each had two sub-sub-helpers? And...

Subdivide the sorting task

H	E	M	G	B	K	A	Q	F	L	P	D	R	C	J	N
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

H	E	M	G	B	K	A	Q
---	---	---	---	---	---	---	---

F	L	P	D	R	C	J	N
---	---	---	---	---	---	---	---

Subdivide again



And again



H E M G

B K A Q

F L P D

R C J N

H E

M G

B K

A Q

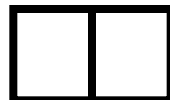
F L

P D

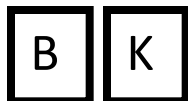
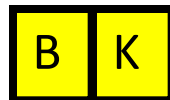
R C

J N

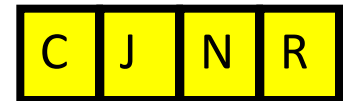
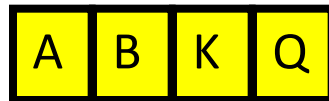
And one last time



Now merge



And merge again



And again



And one last time

A	B	C	D	E	F	G	H	J	K	L	M	N	P	Q	R
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

A	B	E	G	H	K	M	Q
---	---	---	---	---	---	---	---

C	D	F	J	L	N	P	R
---	---	---	---	---	---	---	---

Done!

A	B	C	D	E	F	G	H	J	K	L	M	N	P	Q	R
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---


```
function y = mergeSort(x)
# x is a vector.  y is a vector
# consisting of the values in x
# sorted from smallest to largest.
```

```
n = length(x);
```

```
if n==1
```

```
    y = x;
```

```
else
```

```
    m = floor(n/2);
```

```
    yL = mergeSortL(x(1:m));
```

```
    yR = mergeSortR(x(m+1:n));
```

```
    y = merge(yL, yR);
```

```
end
```

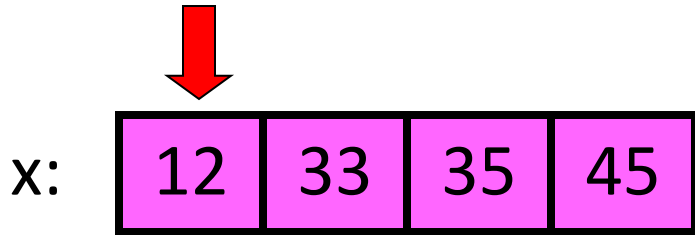
The central sub-problem is the **merging** of two sorted arrays into one single sorted array

12	33	35	45
----	----	----	----

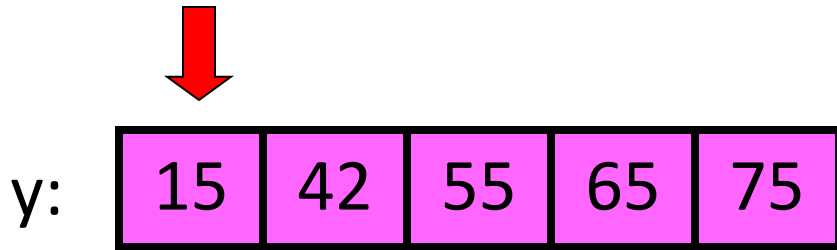
15	42	55	65	75
----	----	----	----	----

12	15	33	35	42	45	55	65	75
----	----	----	----	----	----	----	----	----

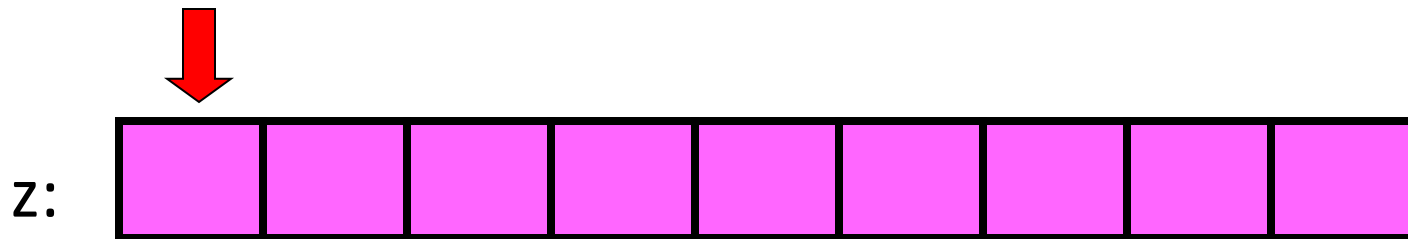
Merge



ix: [1]



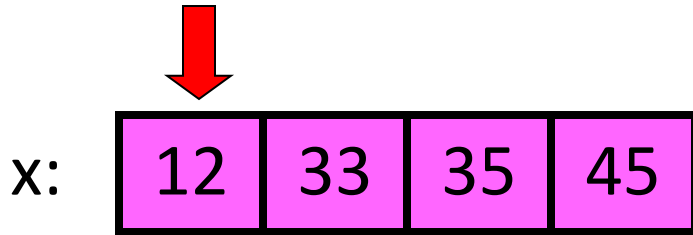
iy: [1]



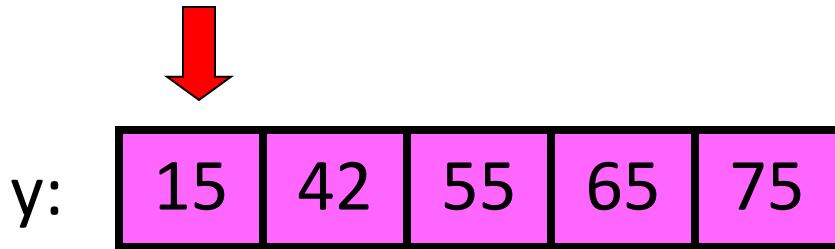
iz: [1]

$ix \leq 4$ and $iy \leq 5$: $x[ix] \leq y[iy]$???

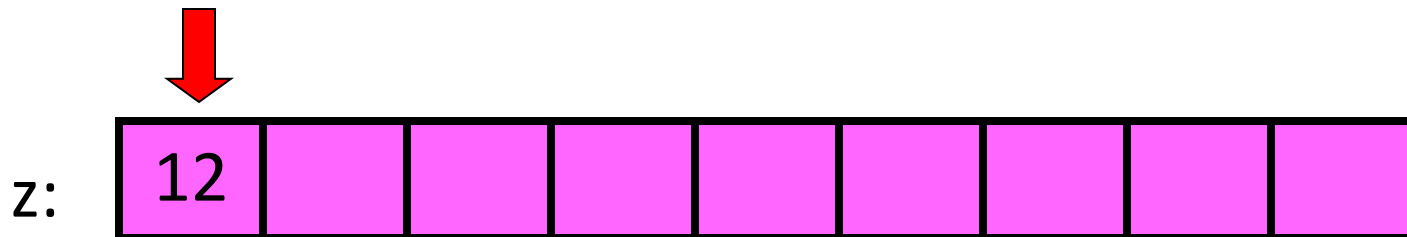
Merge



ix: 



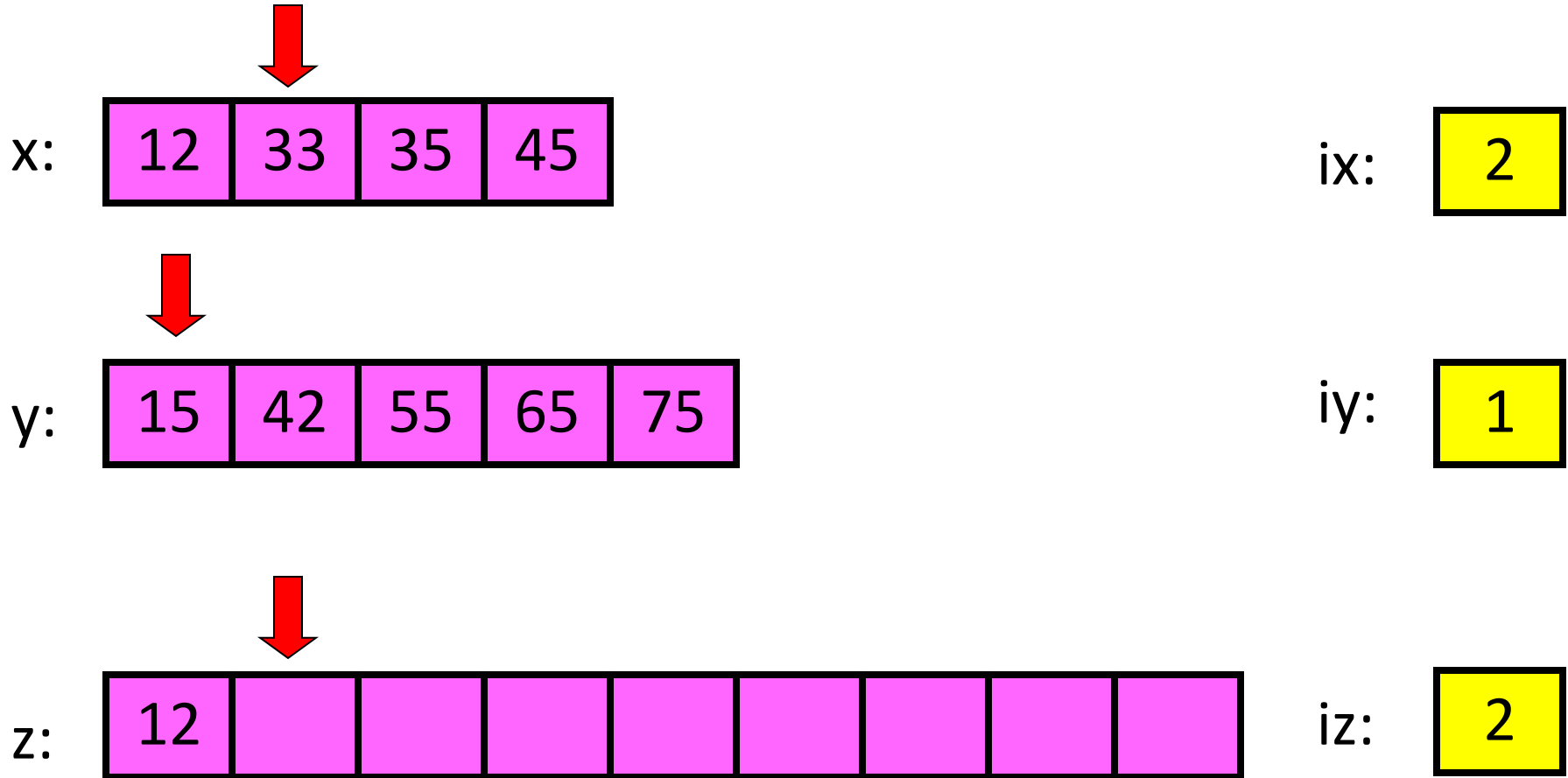
iy: 



iz: 

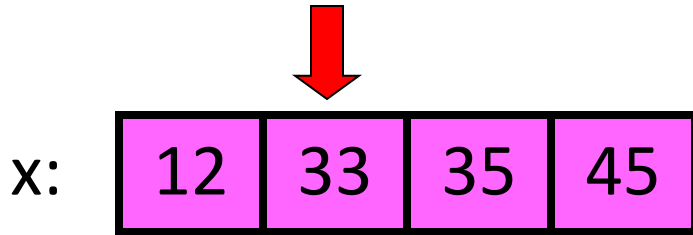
$ix \leq 4$ and $iy \leq 5$: $x(ix) \leq y(iy)$ YES

Merge



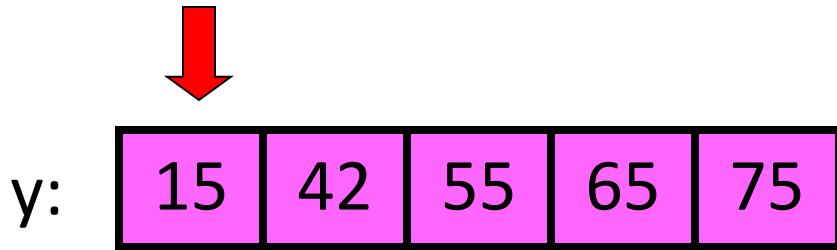
$ix \leq 4$ and $iy \leq 5$: $x[ix] \leq y[iy]$???

Merge



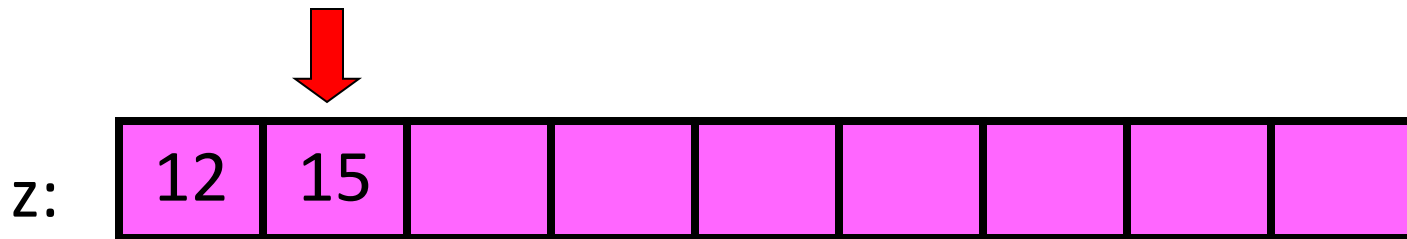
ix:

2



iy:

1

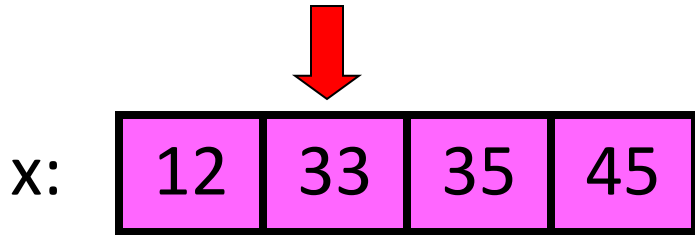


iz:

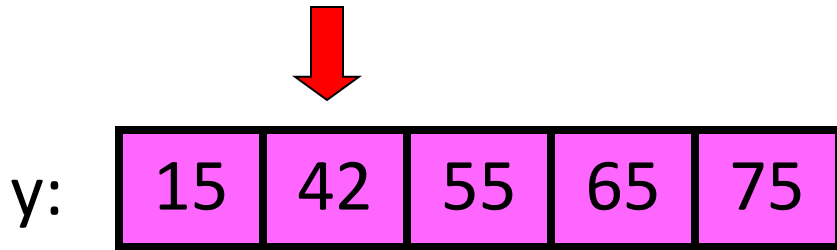
2

$ix \leq 4$ and $iy \leq 5$: $x(ix) \leq y(iy)$ **NO**

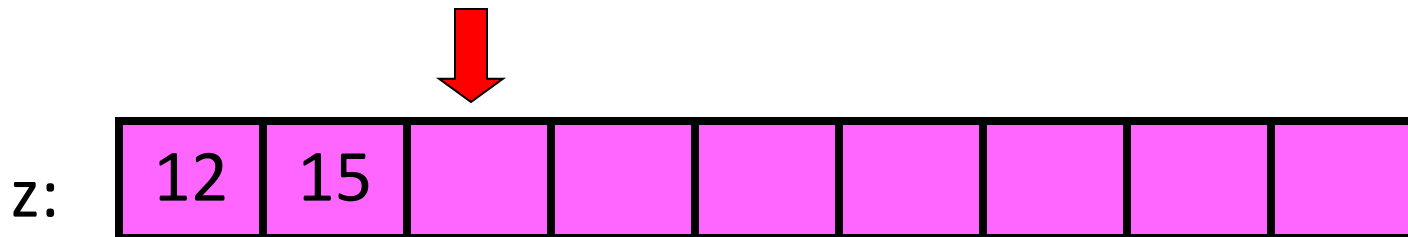
Merge



ix: [2]



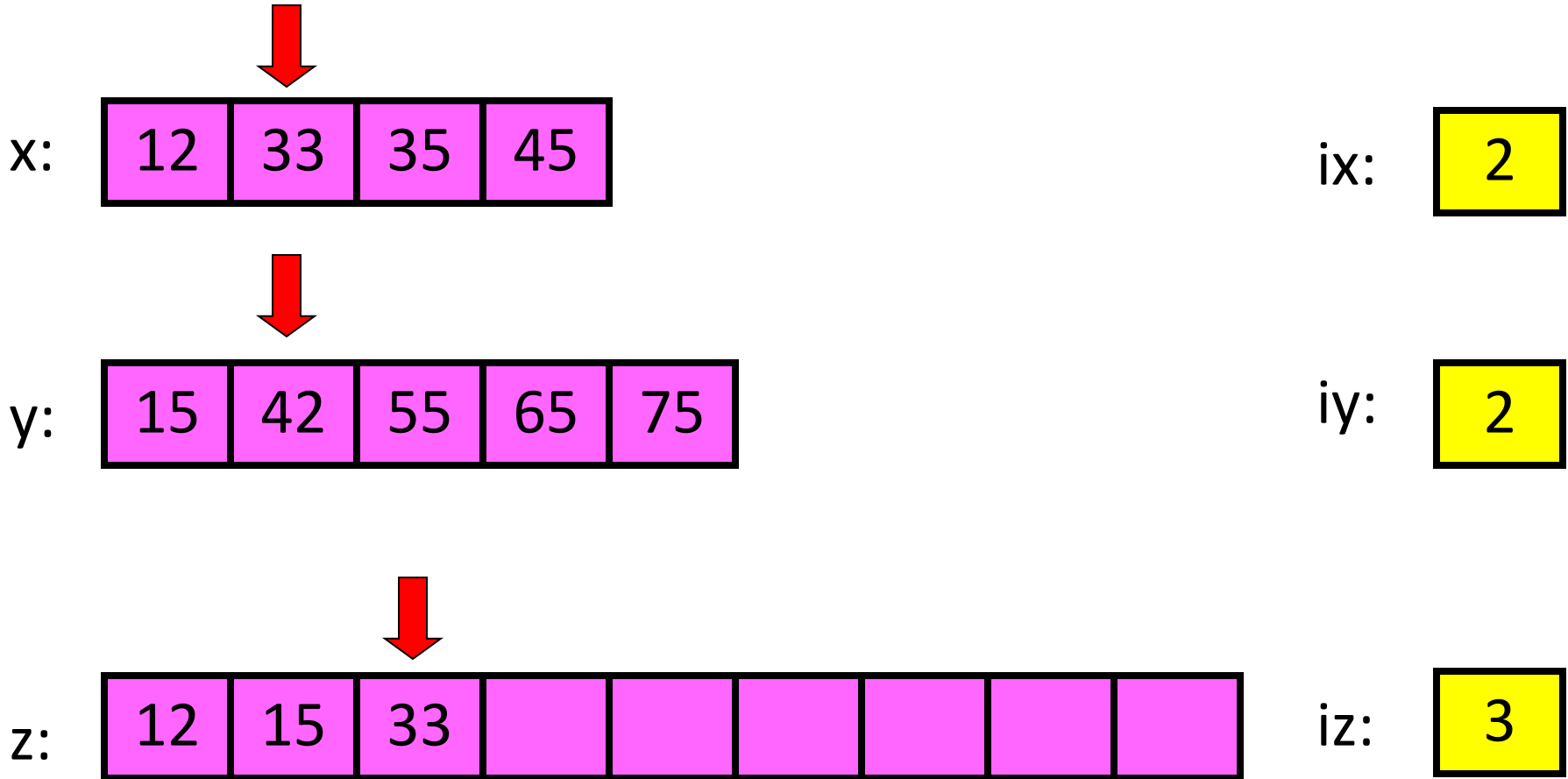
iy: [2]



iz: [3]

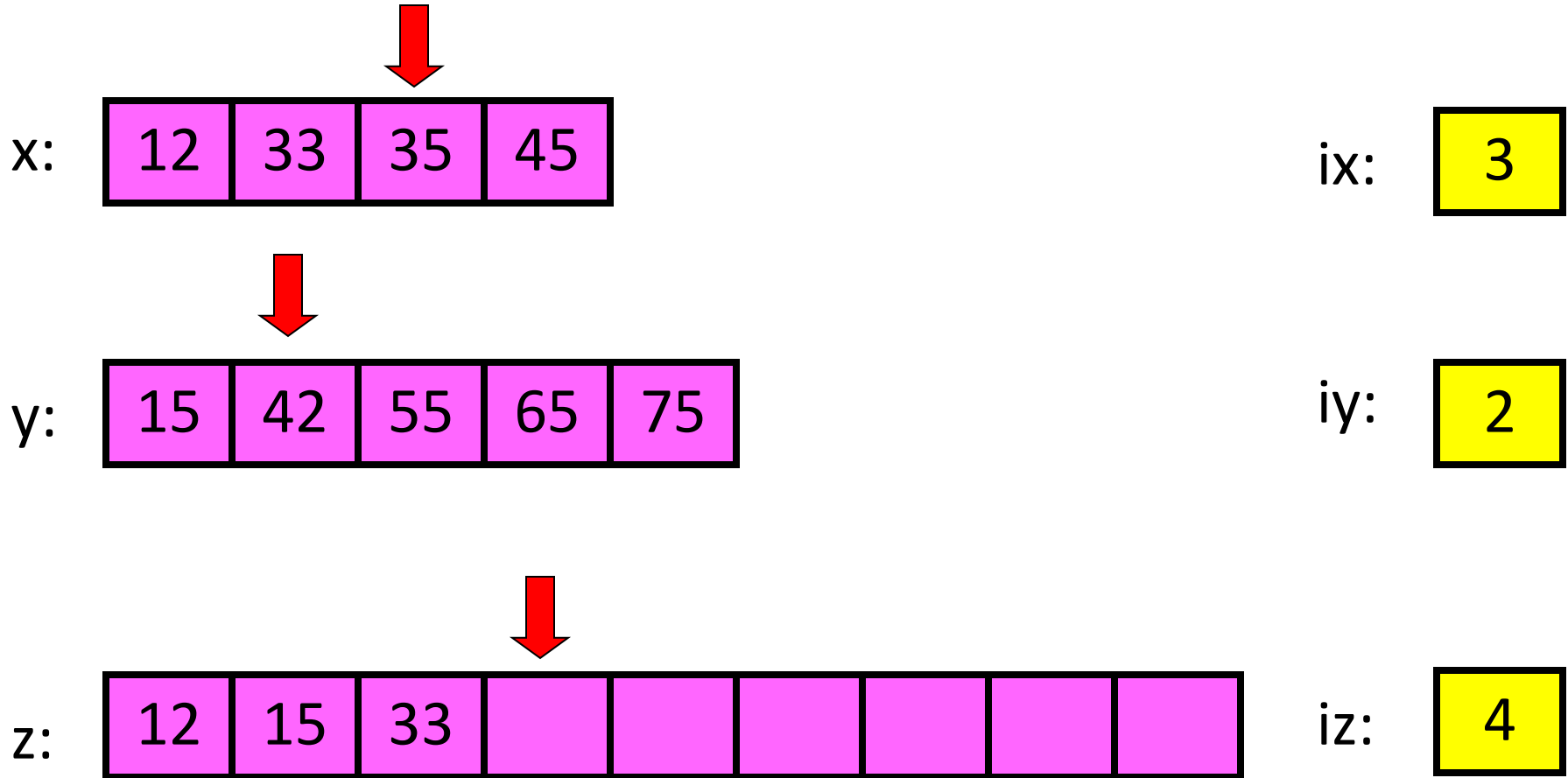
$ix \leq 4$ and $iy \leq 5$: $x[ix] \leq y[iy]$???

Merge



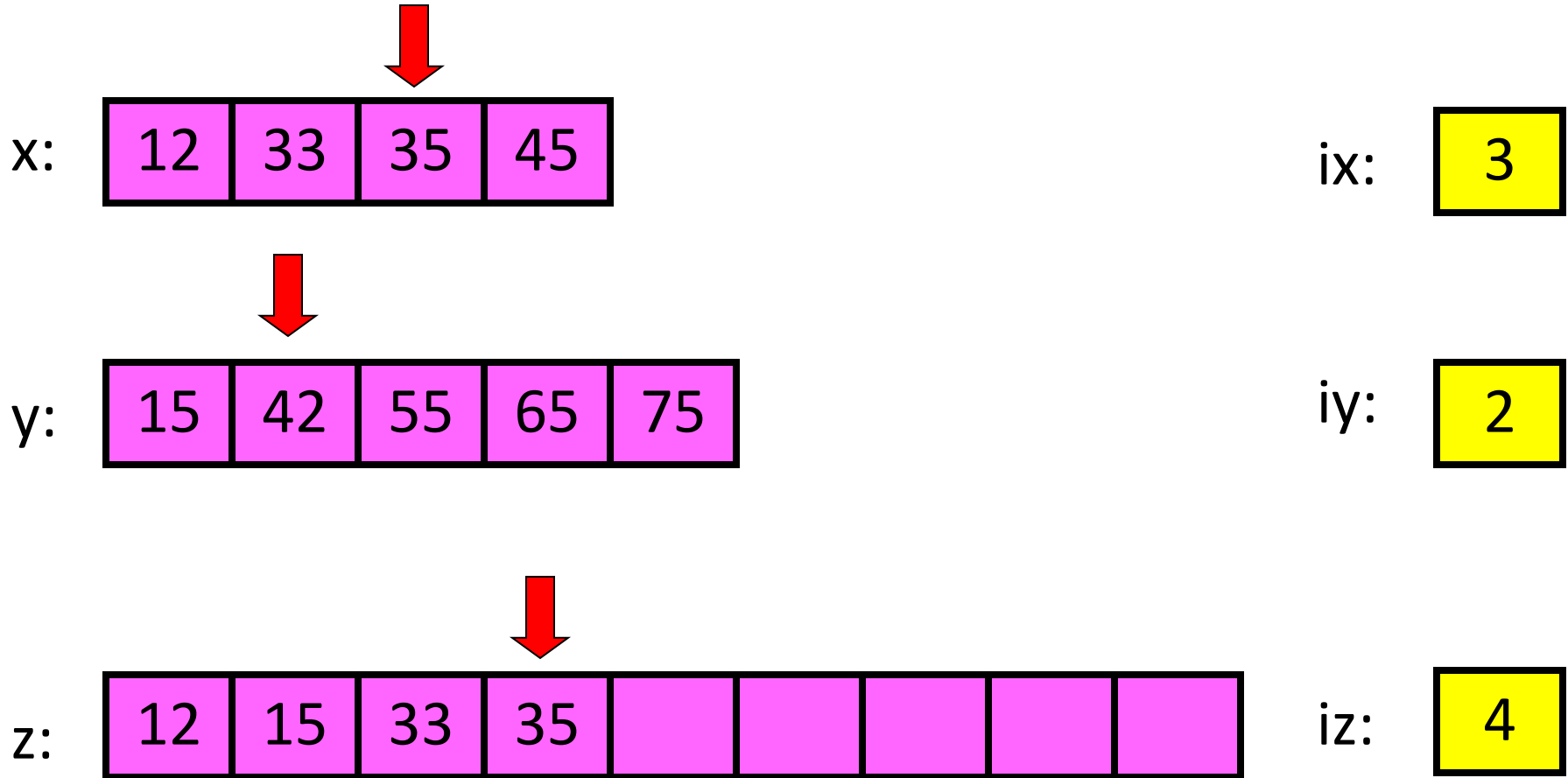
$ix \leq 4$ and $iy \leq 5$: $x(ix) \leq y(iy)$ YES

Merge



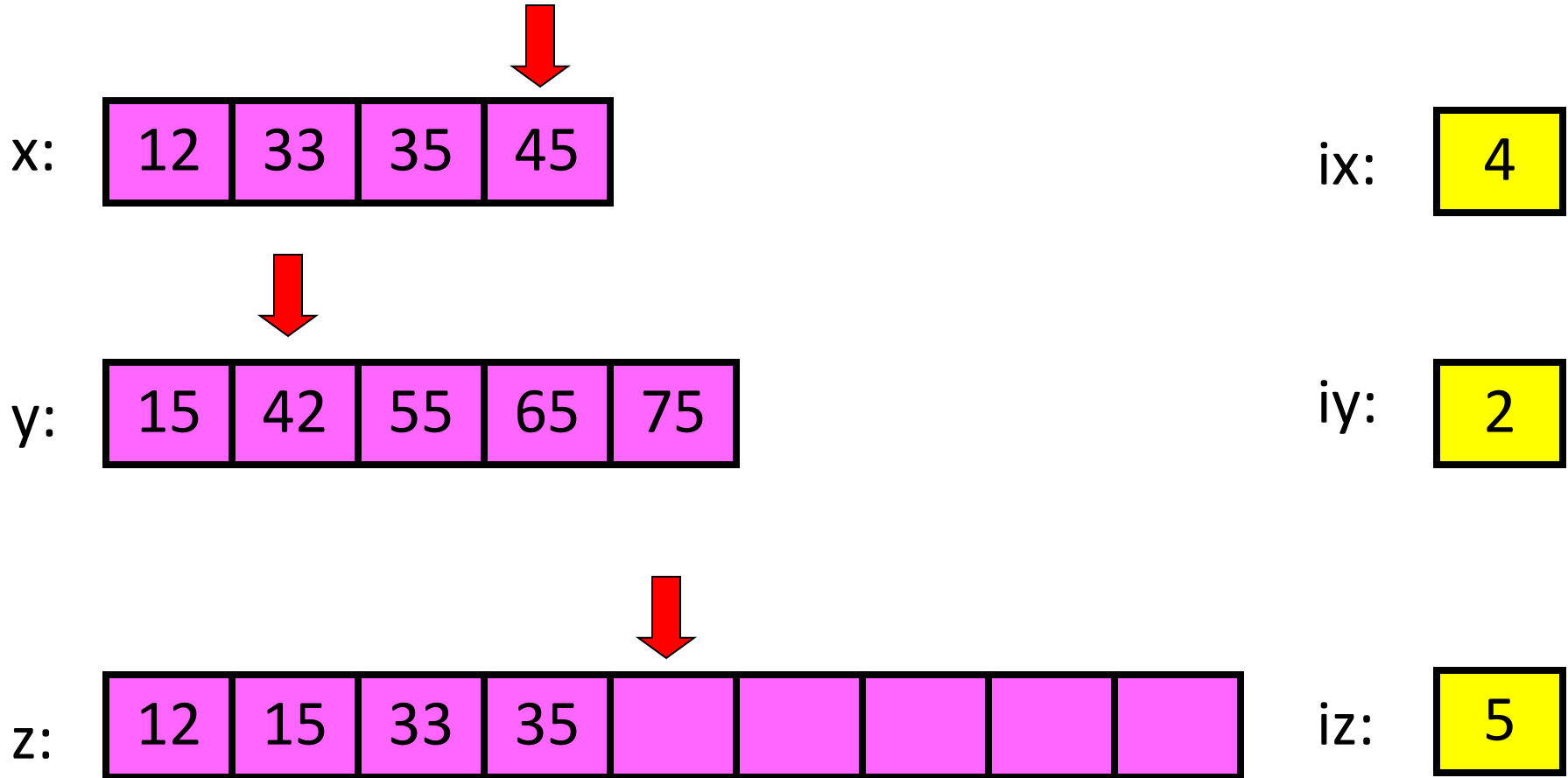
$ix \leq 4$ and $iy \leq 5$: $x(ix) \leq y(iy)$???

Merge



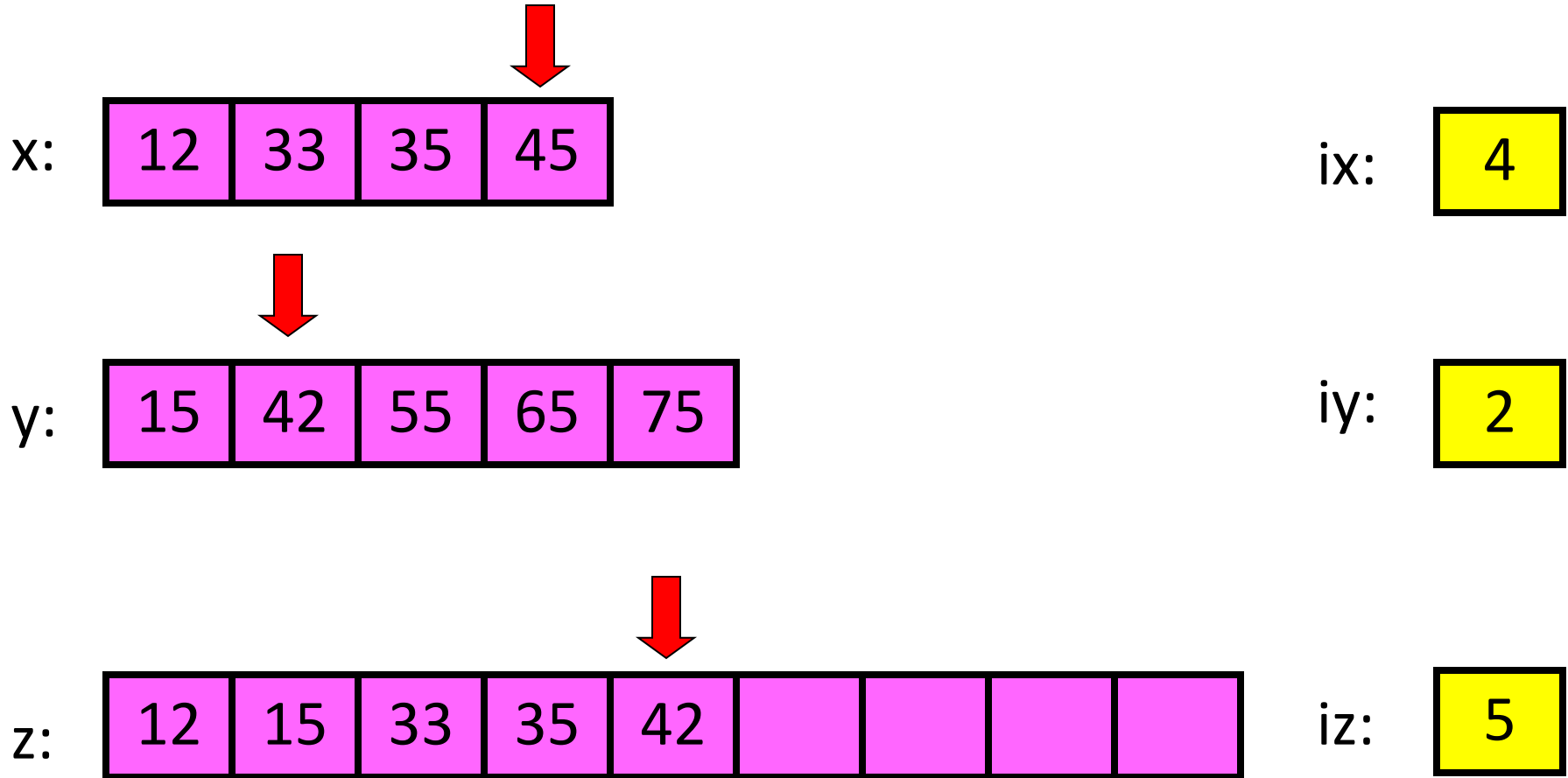
$ix \leq 4$ and $iy \leq 5$: $x(ix) \leq y(iy)$ YES

Merge



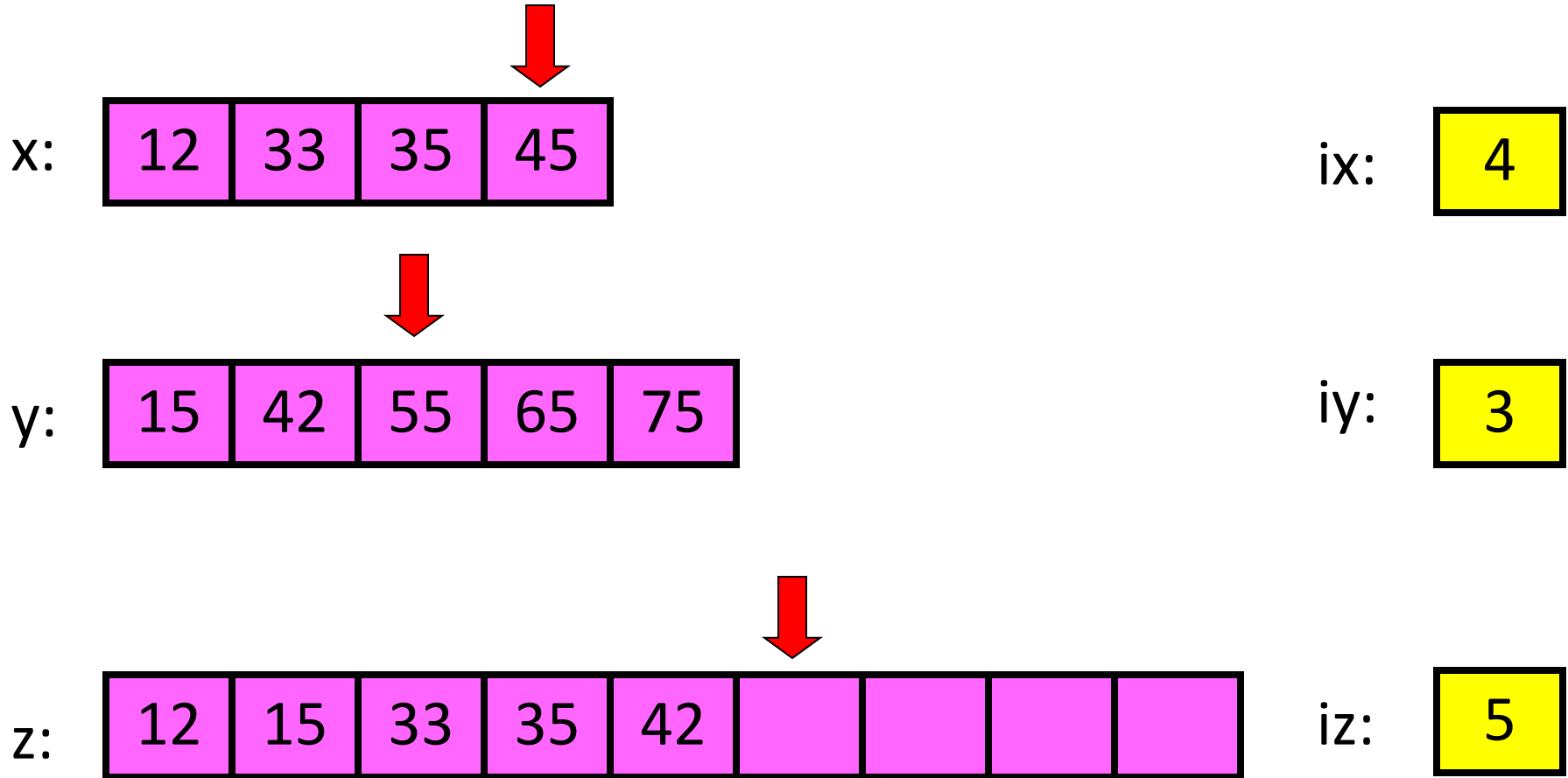
$ix \leq 4$ and $iy \leq 5$: $x(ix) \leq y(iy)$???

Merge



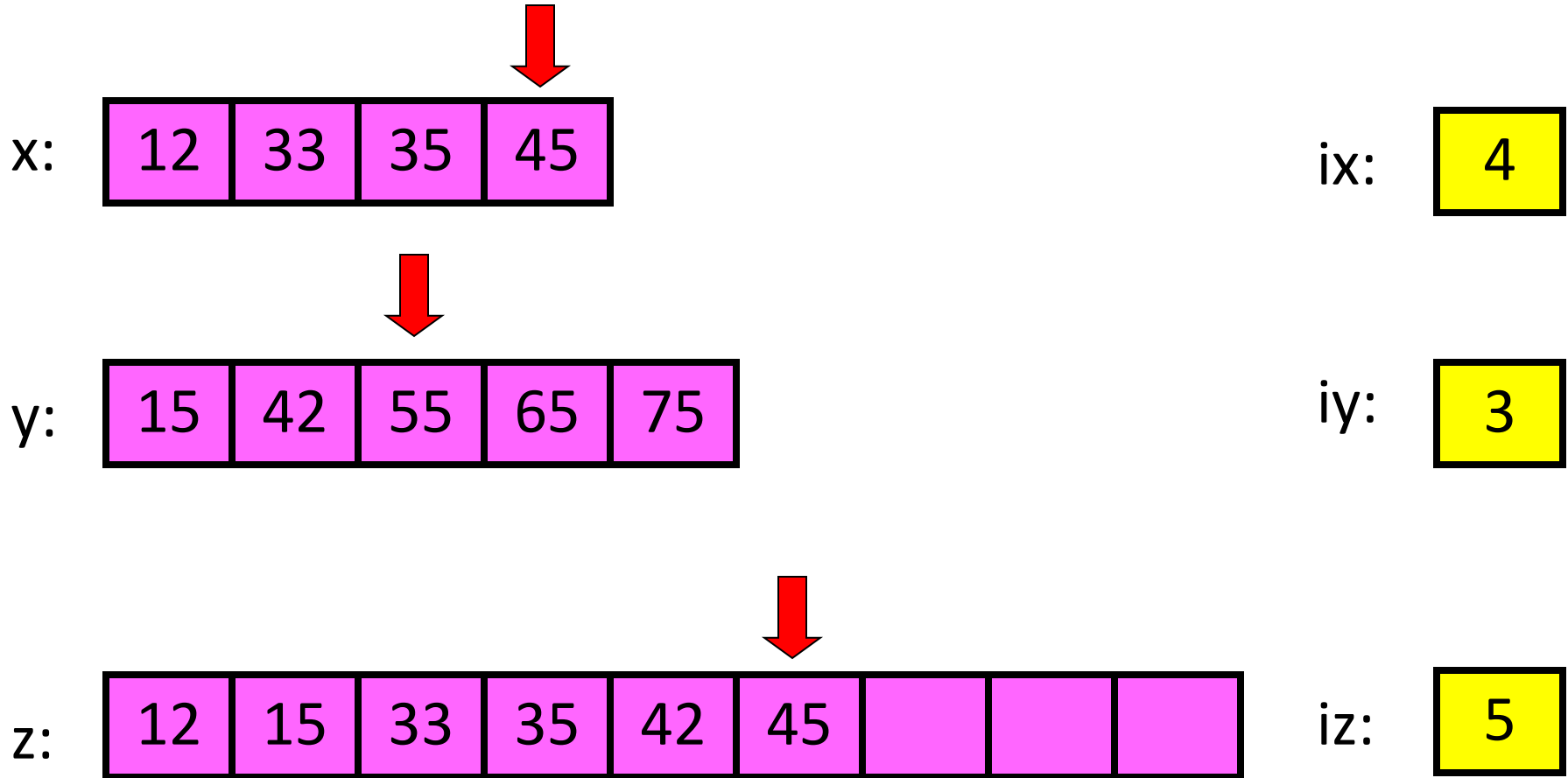
$ix \leq 4$ and $iy \leq 5$: $x(ix) \leq y(iy)$ NO

Merge



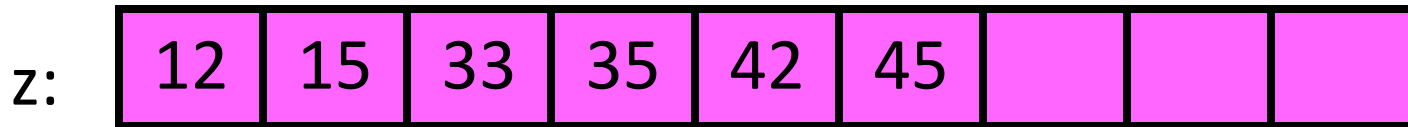
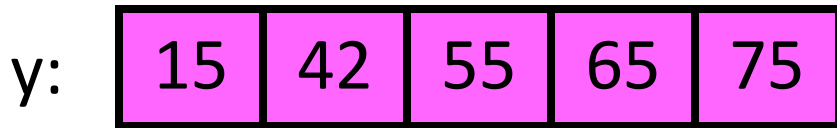
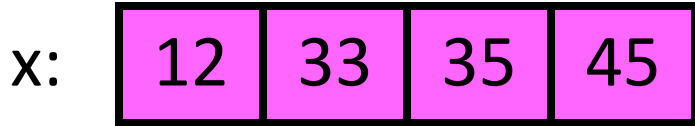
$ix \leq 4$ and $iy \leq 5$: $x(ix) \leq y(iy)$???

Merge



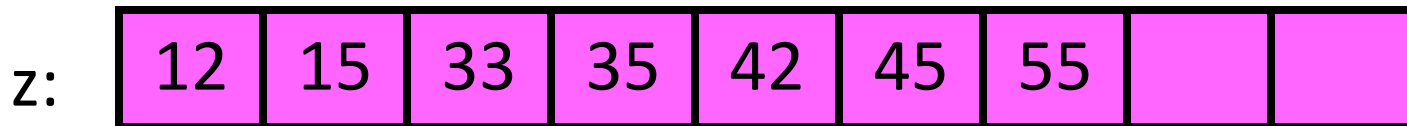
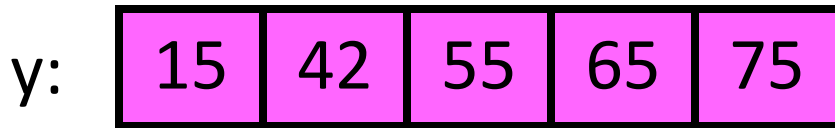
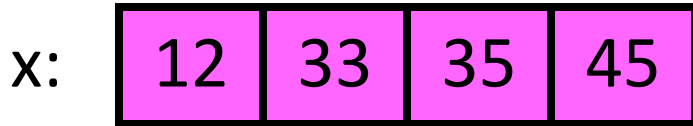
$ix \leq 4$ and $iy \leq 5$: $x(ix) \leq y(iy)$ YES

Merge



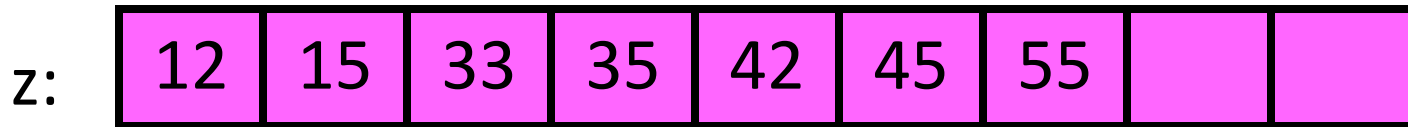
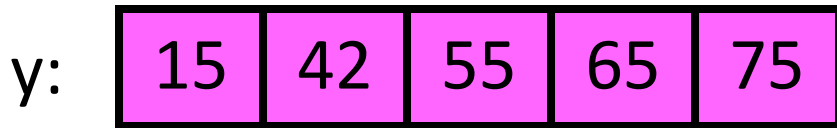
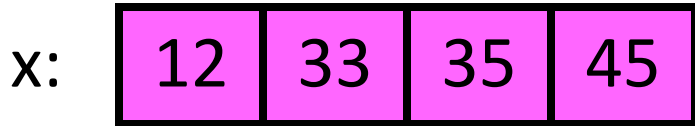
$ix > 4$

Merge



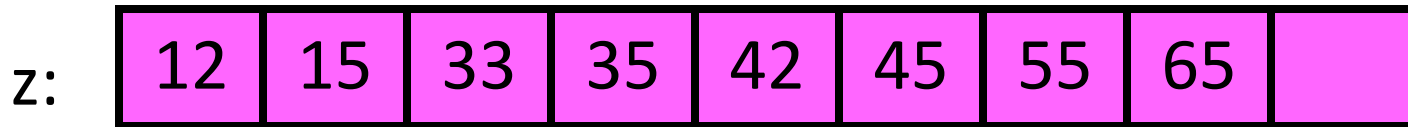
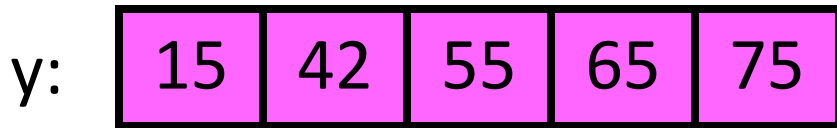
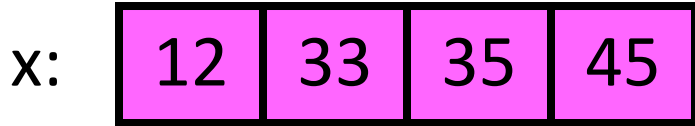
$ix > 4$: take $y(iy)$

Merge



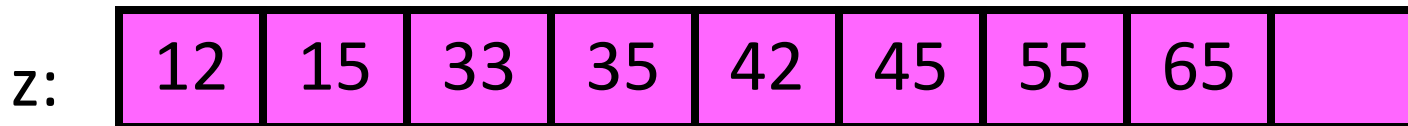
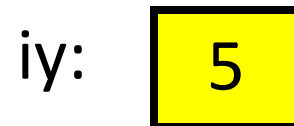
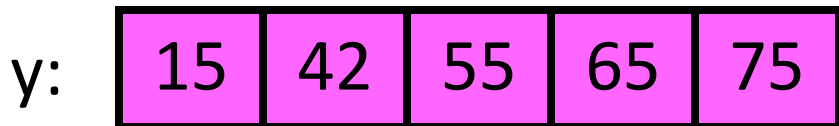
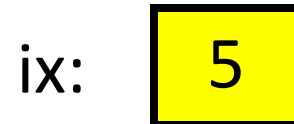
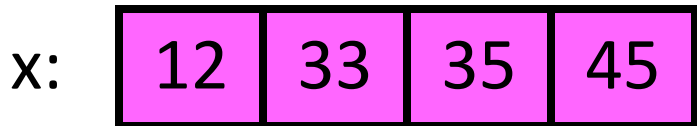
$iy \leq 5$

Merge



$iy \leq 5$

Merge



$iy \leq 5$

Merge



x:

12	33	35	45
----	----	----	----

ix:

5

y:

15	42	55	65	75
----	----	----	----	----

iy:

5

z:

12	15	33	35	42	45	55	65	75
----	----	----	----	----	----	----	----	----

iz:

9

$iy \leq 5$

```
function z = merge(x,y)
nx = length(x); ny = length(y);
z = zeros(1,nx+ny);
ix = 1; iy = 1; iz = 1;
```

```
function z = merge(x,y)
nx = length(x); ny = length(y);
z = zeros(1, nx+ny);
ix = 1; iy = 1; iz = 1;
while ix<=nx && iy<=ny

end
# Deal with remaining values in x or y
```

```
function z = merge(x,y)
nx = length(x); ny = length(y);
z = zeros(1, nx+ny);
ix = 1; iy = 1; iz = 1;
while ix<=nx && iy<=ny
    if x(ix) <= y(iy)
        z(iz) = x(ix); ix=ix+1; iz=iz+1;
    else
        z(iz) = y(iy); iy=iy+1; iz=iz+1;
    end
end
# Deal with remaining values in x or y
```

```
function z = merge(x,y)
nx = length(x); ny = length(y);
z = zeros(1, nx+ny);
ix = 1; iy = 1; iz = 1;
while ix<=nx && iy<=ny
    if x(ix) <= y(iy)
        z(iz) = x(ix); ix=ix+1; iz=iz+1;
    else
        z(iz) = y(iy); iy=iy+1; iz=iz+1;
    end
end
while ix<=nx # copy remaining x-values
    z(iz) = x(ix); ix=ix+1; iz=iz+1;
end
while iy<=ny # copy remaining y-values
    z(iz) = y(iy); iy=iy+1; iz=iz+1;
end
```



```
function y = mergeSort(x)
# x is a vector.  y is a vector
# consisting of the values in x
# sorted from smallest to largest.
```

```
n = length(x);
```

```
if n==1
```

```
    y = x;
```

```
else
```

```
    m = floor(n/2);
```

```
    yL = mergeSortL(x(1:m));
```

```
    yR = mergeSortR(x(m+1:n));
```

```
    y = merge(yL, yR);
```

```
end
```

```
function y = mergeSortL(x)
# x is a vector.  y is a vector
# consisting of the values in x
# sorted from smallest to largest.

n = length(x);
if n==1
    y = x;
else
    m = floor(n/2);
    yL = mergeSortL_L(x(1:m));
    yR = mergeSortL_R(x(m+1:n));
    y = merge(yL, yR);
end
```

```
function y = mergeSortL_L(x)
# x is a vector.  y is a vector
# consisting of the values in x
# sorted from smallest to largest.
```

```
n = length(x);
```

```
if n==1
```

```
    y = x;
```

```
else
```

```
    m = floor(n/2);
```

```
    yL = mergeSortL_L_L(x(1:m));
```

```
    yR = mergeSortL_L_R(x(m+1:n));
```

```
    y = merge(yL, yR);
```

```
end
```

There should be just one mergeSort function!

```
function y = mergeSort(x)
# x is a vector.  y is a vector
# consisting of the values in x
# sorted from smallest to largest.
```

```
n = length(x);
```

```
if n==1
```

```
    y = x;
```

```
else
```

```
    m = floor(n/2);
```

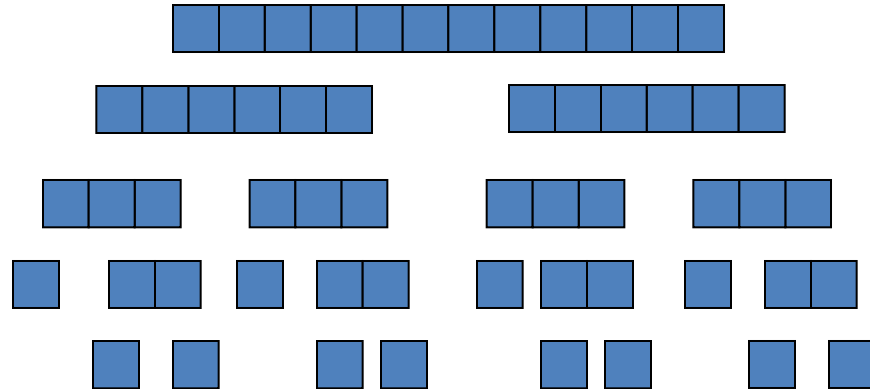
```
    yL = mergeSort(x(1:m));
```

```
    yR = mergeSort(x(m+1:n));
```

```
    y = merge(yL, yR);
```

```
end
```

```
function y=mergeSort(x)
n=length(x);
if n==1
    y=x;
else
    m=floor(n/2);
    yL=mergeSort(x(1:m));
    yR=mergeSort(x(m+1:n));
    y=merge(yL,yR);
end
```



Divide-and-Conquer Examples

- Sorting: mergesort and quicksort
- Binary tree traversals
- Binary search
- Multiplication of large integers
- Matrix multiplication: Strassen's algorithm
- Closest-pair algorithm