

198:314, Sects1-3 Principles of Programming Languages Fall 2001

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Introduction

- **Administrivia**
- **Why study PLs?**
- **Paradigms: imperative, functional, object-oriented, logic**
- **History of PLs**

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- **Class webpage**
 - <http://remus.rutgers.edu/cs314>
 - Look at for course information, rules, grading, important dates, etc.
 - Read policy on academic dishonesty
 - Co-ordinated class, all programming projects, book homeworks and tests are same across all lectures
- **MW5 Lecture webpage**
 - <http://remus.rutgers.edu/cs314/f2001/ryder/>
 - Lecture notes available online in pdf
 - Only print out 2-up or 4-up to save paper
 - Recitation attendance and participation *counts* in your final grade
- **Four programming projects posted and submitted for grading electronically**

Syllabus

- **Introduction**
- **Formal Languages - RE's, FSA's**
- **Logic Programming (Prolog)**
- **Names and Binding**
- **Imperative Programming (C)**
- **Block Structure**
- **Object-oriented Programming (C++)**
- **Types**
- **Functional Programming (Scheme)**
- **Formal Languages - Grammars**
- **Concurrency**

Course Goals

- **To gain understanding of basic structures of programming languages**
 - Types, control structures, naming conventions
- **To study different language paradigms**
 - To ensure an appropriate language is selected for a task
 - Object-oriented, functional, imperative, logic
- **To make learning new programming languages easier by knowing shared features**

What is a programming language?

“a language intended for use by a **person** to express a **process** by which a **computer** can solve a problem” -Hope and Jipping

“a set of conventions for communicating an algorithm” - E. Horowitz

“ the art of programming is the art of organizing complexity” - Dijkstra, 1972

Why learn more than one PL?

- **Each language paradigm encourages thinking about a problem in a particular manner**
 - Finding a natural match between problem and PL
- **Somewhat different functionality supplied by different paradigms**
- **Computer professionals must be multi-lingual**
 - PLs change over time as computer architecture changes
 - Specific applications sometimes result in specialized PLs
 - Need to understand each PLs functionality and limitations

Imperative Paradigm

- **Underlying notion of an abstract machine**
 - **Von Neumann architecture**
 - Store (memory)
 - Accumulator (ALU)
 - Load/store into memory
 - **Key operation: assignment**

Imperative PLs

Sum up twice each
number from 1 to N.

Fortran

```
SUM = 0
DO 11 K=1, N
SUM = SUM + 2*K
11 CONTINUE
```

C

```
sum = 0;
for (k = 1; k <= n; ++k)
    sum += 2*k;
```

Pascal

```
sum := 0;
for k := 1 to n do
    sum := sum + 2*k;
```

Functional Paradigm

- **Process of problem solution expressed as a sequence of operations on the data**
 - (Pure) value binding through parameter passing
 - No store accessible through names
 - No iteration
 - Key operation: function application (with recursion)

Functional Paradigm

Scheme

```
(define (sum n)
  (if (= n 0)
      0
      (+ (* n 2) (sum (- n 1))))
)

(sum 4) evaluates to 20
```

Logic Paradigm

- **Program is a formal description of characteristics required of a problem solution**
 - Programs tell *what should be not how to make it so*
 - Solutions through a reasoning process called **theorem proving**
 - **Key operation: unification**

Logic Paradigm

```
sum(0, 0).  
sum(N, S) :- NN is N - 1,  
            sum(NN, SS),  
            S is N * 2 + SS.
```

Prolog

```
?- sum(1, 2).  
yes  
?- sum (2, 4).  
no  
?- sum(20, S).  
S = 420  
?- sum (X, Y).  
X = 0 = Y
```

Object-oriented Paradigm

- **Organizes a program to be operations on abstract representations of the data**
 - **Objects with data abstraction and information hiding**
 - Object implementation is hidden from user
 - **Actions performed on objects (messages)**
 - **Can combine with imperative or functional paradigm easily**
 - **Key operation: message passing**

Object-oriented Paradigm

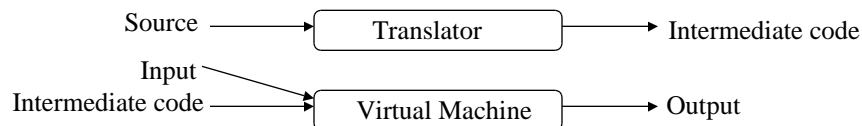
```
class intSet : public Set
{ public: intSet () { }
//inherits Set add_element(), Set del_element()
//from Set class, defined as a set of Objects
public int sum() {
    int s = 0;
    SetEnumeration e = new SetEnumeration (this);
    while (e.hasMoreElements ()) do
    { s =s + ((Integer)e. nextElement ().intValue ()); }
    return s;
}
}
```

Java

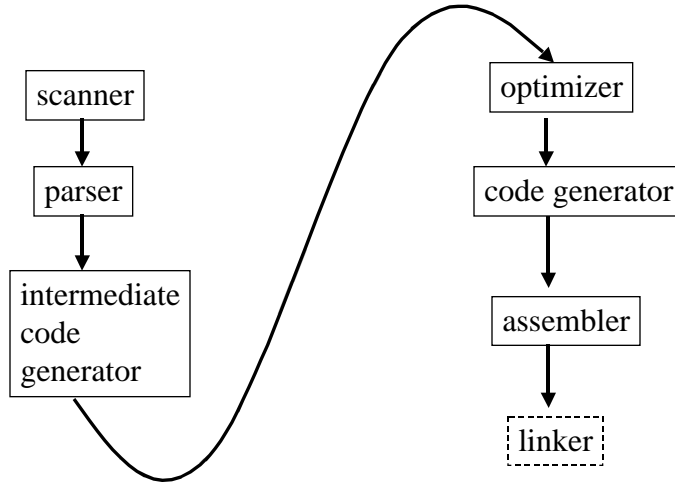
Translation

- **Compilation:** translation of a program written in a high-level PL into a form that is executable on the machine
- **Interpretation:** a program is translated and executed one statement at a time
- Most PL systems are a mixture of these two
 - **Interpreted:** Java, Scheme, Prolog
 - **Compiled:** Fortran, C, C++

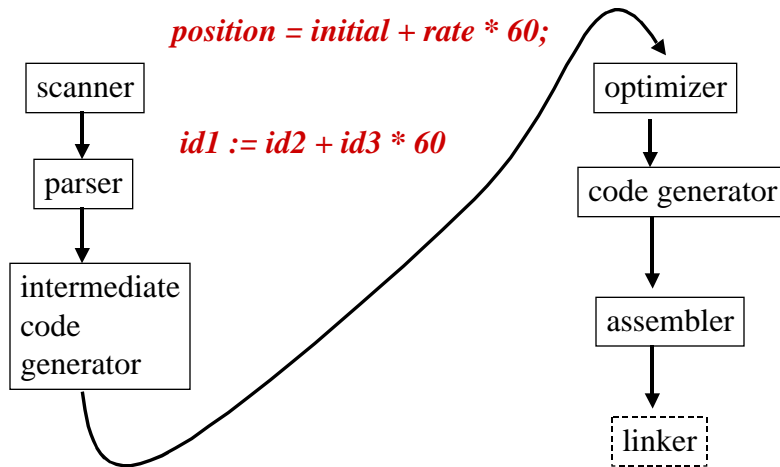
Scott, p 10



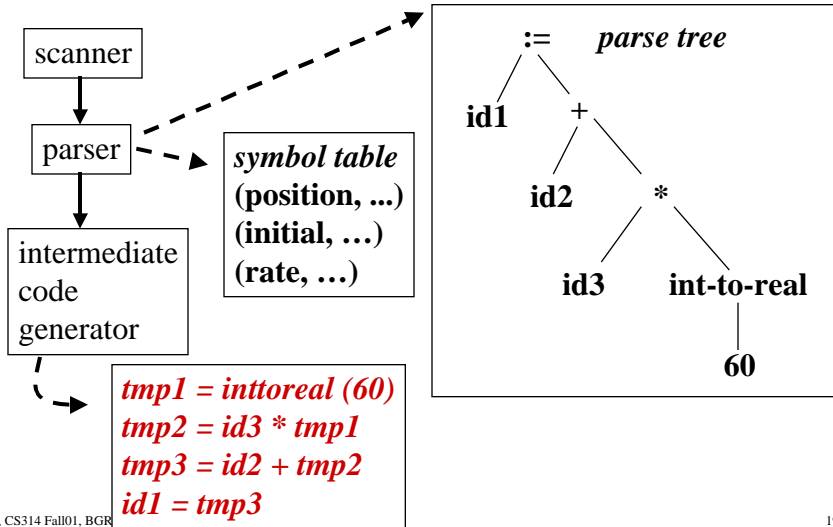
Compilation



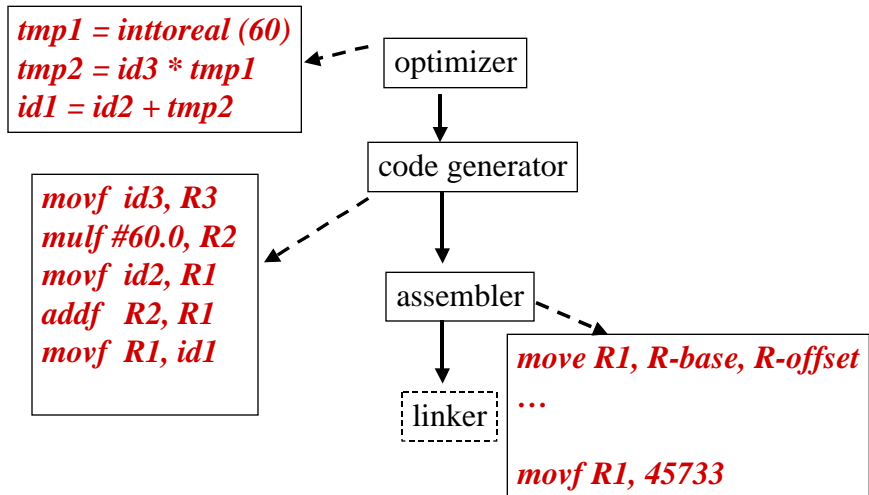
Compilation



Compilation



Compilation



Desiderata for PL Design

- **Readability**
 - comments, names, (...) syntax
- **Simple to learn**
- **Orthogonal**
 - small number of concepts combine regularly and systematically (without exceptions)
- **Portability**
 - language standardization
- **Abstraction**
 - data and control abstractions

History of PLs

- **Prehistory**
 - 300 B.C. Greece, Euclid invented the greatest common divisor algorithm - *oldest known algorithm*
 - ~1820-1850 England, Charles Babbage invented 2 mechanical computational devices
 - difference engine
 - analytical engine
 - Countess Ada Augusta of Lovelace, *first computer programmer*
 - **Precursors to modern machines**
 - 1940's United States, ENIAC developed to calculate trajectories

History of PLs - 2

- **1950's United States, first high-level PLs invented**
 - **Fortran** 1954-57, John Backus (IBM on 704) designed for numerical scientific computation
 - fixed format for punched cards
 - implicit typing
 - only counting loops, if test versus zero
 - only numerical data
 - 1957 optimizing Fortran compiler translates into code as efficient as hand-coded

History of PLs - 3

- **Algol60** 1958-60, designed by international committee for numerical scientific computation [Fortran]
 - block structure with lexical scope
 - free format, reserved words
 - while loops, recursion
 - explicit types
 - BNF developed for formal syntax definition
- **Cobol** 1959-60, designed by committee in US, manufacturers and DoD for business data processing
 - records
 - focus on file handling
 - English-like syntax

History of PLs - 4

- ***APL*** 1956-60 Ken Iverson, (IBM on 360, Harvard) designed for array processing
 - functional programming style
- ***LISP*** 1956-62, John McCarthy (MIT on IBM704, Stanford) designed for non-numerical computation
 - uniform notation for program and data
 - new conditional control structure (COND)
 - recursion as main control structure
- ***Snobol*** 1962-66, Farber, Griswold, Polansky (Bell Labs) designed for string processing
 - powerful pattern matching

History of PLs - 5

- ***PL/I*** 1963-66, IBM designed for general purpose computing [Fortran, Algol60, Cobol]
 - user controlled exceptions
 - multi-tasking
- ***Simula67*** 1967, Dahl and Nygaard (Norway) designed as a simulation language [Algol60]
 - data abstraction
 - inheritance of properties
- ***Algol68*** 1963-68, designed for general purpose computing [Algol60]
 - orthogonal language design
 - interesting user defined types

History of PLs - 6

- **Pascal** 1969, N. Wirth(ETH) designed for teaching programming [Algol60]
 - 1 pass compiler
 - call-by-value semantics
- **Prolog** 1972, Colmerauer and Kowalski designed for Artificial Intelligence applications
 - theorem proving with unification as basic operation
 - logic programming
- Recent
 - **C** 1974, D. Ritchie (Bell Labs) designed for systems programming
 - allows access to machine level within high-level PL
 - efficient code generated

History of PLs - 7

- **Clu** 1974-77, B. Liskov (MIT) designed for simulation [Simula]
 - supports data abstraction and exceptions
 - precise algebraic language semantics
 - attempt to enable verification of programs
- **Smalltalk** mid-1970s, Alan Kay (Xerox Parc), considered first real object-oriented PL, [Simula]
 - encapsulation, inheritance
 - easy to prototype applications
 - hides details of underlying machine
- **Scheme** mid-1970s, Guy Steele, Gerald Sussman (MIT)
 - Static scoping and first-class functions

History of PLs - 8

- **Concurrent Pascal** 1976 Per Brinch Hansen (U Syracuse) designed for asynchronous concurrent processing [Pascal]
 - monitors for safe data sharing
- **Modula** 1977 N. Wirth (ETH), designed language for large software development [Pascal]
 - to control interfaces between sets of procedures or *modules*
 - real-time programming
- **Ada** 1979, US DoD committee designed as general purpose PL
 - explicit parallelism - *rendezvous*
 - exception handling and generics (*packages*)

History of PLs - 9

- **C++** 1985, Bjarne Stroustrup (Bell Labs) general purpose
 - goal of type-safe object-oriented PL
 - compile-time type checking
 - templates
- **Java** ~1995, J. Gosling (SUN)
 - aimed at portability across platform through use of JVM - abstract machine to implement the PL
 - aimed to *fix* some problems with previous OOPLs
 - multiple inheritance
 - static and dynamic objects
 - ubiquitous exceptions
 - thread objects