

# Method Resolution Approaches

- **Static - procedural languages (w/o fcn ptrs)**
- **Dynamically determined by data values**
  - C with function pointers
  - Compile-time analysis can estimate possible callees
- **Dynamically determined by receiver type**
  - Some polymorphic OOPLs use runtime type of first parameter to specialize behaviors
  - Some OOPLs also use runtime types of other arguments
  - **Problem:** how to have an efficient implementation of this kind of dynamic dispatch?

## Dynamic Dispatch

- **Choices PLs have to make:**
  - When resolve function targets?
  - What to look at to do the method resolution? (e.g., receiver runtime type? Argument runtime types?)
  - How to divide the work between runtime and compile time?
  - Emphasize flexibility or performance?

# Method Redefinition

- **Overriding** - replacing a superclass's implementation of a method, by one with identical signature (except receiver type)
  - Method must be accessible, non-static
- **Overloading** - providing more than one method with same name, but different signatures to distinguish them
- Simple cases of both are intuitive

# Inheritance

- Overriding can widen method visibility
- Can override instance variables, but can still get to superclass variable using *super*
- Preferred inheritance uses all *private* data and provides *observer* and *mutator* methods
  - Using *geta()*, *seta()* methods means that changing superclass structure will not affect subclasses
- Access to
  - Methods is by run-time type of object referenced
  - Instance variables is by compile-time type of reference

# Possible Cases

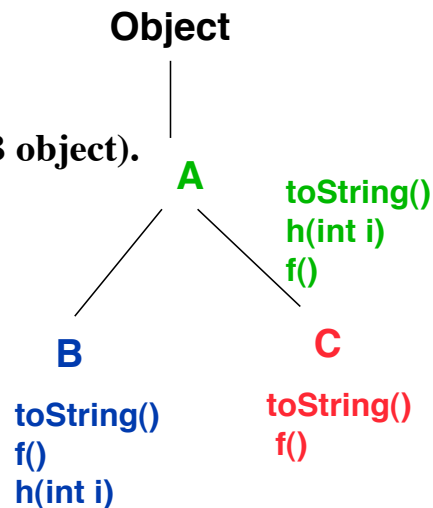
- **Inheritance**, but all method names are unique
- **Inheritance with *overriding***
  - Lookup happens at run-time based only on receiver's class
  - Next slide: A,B,C with respect to f(); A,B wrt h()
- **Inheritance with *overloading* (different method signatures)**
  - Java: Lookup establishes best match type signature at compile-time based on arguments' and receiver's declared classes; actual binding done by run-time lookup to match selection
  - Next slide, A,B wrt s()

# Overriding Example

```
A a = new B();
```

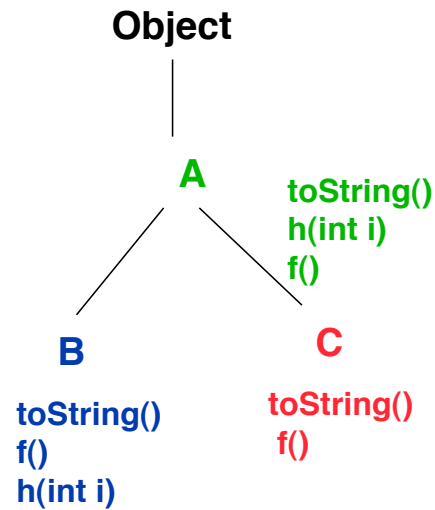
How to resolve **a.toString()**?

1. At **run-time**, determine class of the object (e.g., a refers to an B object).
2. Start lookup for method with same signature in class B.
3. Proceed up inheritance hierarchy until find closest superclass with same signature (i.e., method toString() ); this may be class B itself



# Java Overriding Example

```
//overriding - fcns have
//same signature
A a1 = new B();
A a2 = new C();
B b = new B();
A a = new A();
a.f(); //A's f()
a1.f(); //B's f()
a2.f(); //C's f()
b.h(0); //B's h()
a1.h(2); //B's h()
a2.h(1); //A's h()
a.h(3); //A's h()
```



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# C++ Approach to Overriding

- If return type and signature of 2 functions match exactly, the 2nd is a redeclaration of the first and is an override
- If signatures of 2 functions match exactly, but return types differ, then 2nd declaration is in *error*
- If signatures differ in number or type of arguments, the 2 function instances are **OVERLOADED**. (return type not considered as part of signature here)

S. Lippman,  
C++ Primer

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# Overloading

- Java chooses to optimize dynamic dispatch by partially resolving references through preprocessing at compile-time
  - Need to use **declared type and number of arguments + receiver type** to help select an unique method
- Results in a not-just-dynamic lookup procedure because pre-selection is done
  - Different from multi-methods (e.g., in Cecil) where dynamic lookup is based on *run-time* types of **receiver** and the **arguments!**

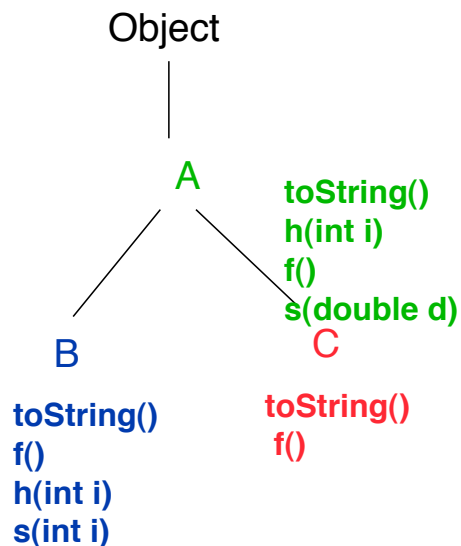
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## Example

```
//overloading -when signatures
//not same, must look at type
//matching between arg and
//param
a.s(3.); //A's s()
a1.s(3.); //A's s() because
//arg is a double and B's
//s() expects an int
b.s(0); //B's s()
b.s(1.0); //A's s()
//casting is not type
//conversion in Java
((A) a1).h(4); //uses B's h()
//matching rules are not
//always straight-forward
a1.s(0); //A's s()
```

```
A a1 = new B();
A a2 = new C();
B b = new B();
A a = new A();
```



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# Overloading Resolution in Java

- **At compile-time**, assemble a set of methods whose parameters are type compatible with the arguments and receiver
- For each invocation
  - Look at compile-time class of receiver and arguments
  - Move up class hierarchy from declared receiver type class trying for a match (possibly widening argument or receiver types)
  - Collect all possible matching methods into a set and then find the **most specific match** (defined on next slide)

## Most Specific Match

- If find unique method with exact match in type and number of arguments and compatible receiver type, choose it.
- Otherwise,
  - If any method *f* has arguments + receiver that can be assigned to any other method *g* in the set, discard *g*; Repeat as much as possible.
  - If only 1 method remains, use it as *template*.
  - If more than 1 method remains, the invocation is ambiguous, so the invoking code is invalid. **Compile-time error!!**

# Overloading Resolution in Java

## Run-time Overriding

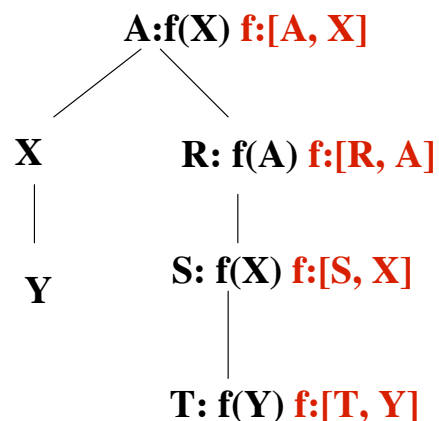
- At run-time, use run-time type of receiver to start search up class hierarchy for function **exactly matching previously defined *template***. (Note: ignore run-time types of arguments)
- Stop going up the hierarchy when find first match to *template* type. Overloading guarantees there will be at least one match.

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## Java Example (cf Don Smith)

- Class hierarchy as shown contains 4 variants of method `f()`
- Signatures[...] include compile-time types of receiver and argument.
- Objects named for their compile-time type

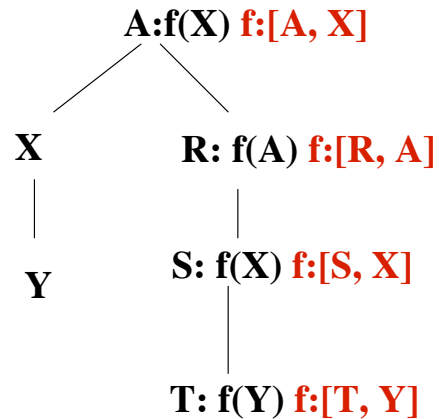


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# Java Example

- **a.f(x)**
  - signature  $f:[A, X]$
  - check classes A
  - matches **f:[A, X]**
- **s.f(a)**
  - signature  $f:[S, A]$
  - check classes S, R, A
  - matches **f:[R, A]**



# Java Example

- **s.f(y)**
  - signature  $f:[S, Y]$
  - checks S, R, A
  - matches **[S, X], [R, A], [A, X]**.

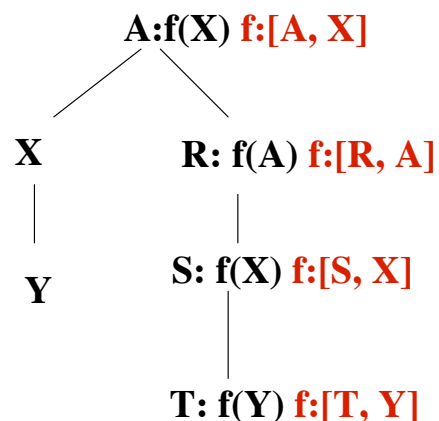
- check pairwise for most specific

[S, X] with [R, A]

[S, X] with [A, X]

[R, A] with [A, X]

**[S, X] is choice**

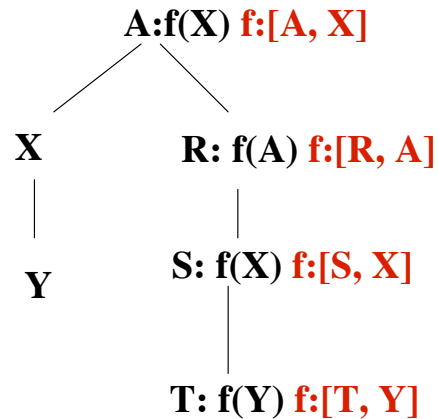


If any method f has arguments + receiver that can be assigned to any other method g in the set, discard g;



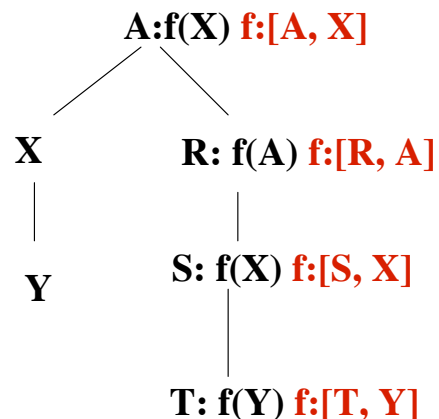
# Java Example

- **r.f(x)**
  - signature  $f:[R, X]$
  - checks **R, A**
  - matches **[R, A], [A, X]**
- **check pair**
  - **R << A but A >> X**
  - **incomparable**
  - **no match**
  - **compile-time ERROR!**



# Java Example

- **t.f(y)**
  - signature  $f:[T, Y]$
  - checks **T, S, R, A**
  - matches **[T, Y], [S, X], [R, A], [A, X]**
- **pairwise check and get [T, Y]**



```
T t = new T();
X x = new Y();
```

...

```
t.f(x);
```

signature is [T, X]

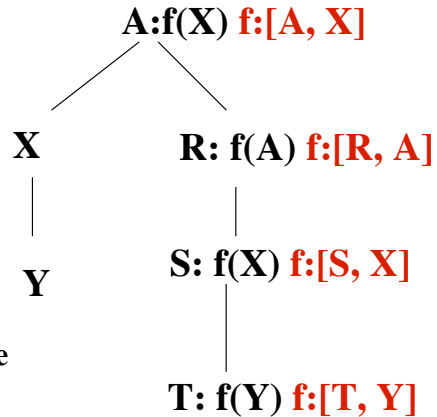
checks T, S, R, A

matches [S,X], [R,A], [A,X]

specificity eliminates all but [S, X]

at run-time receiver is class T and argument is class Y. however, call will be resolved to f(X) in class S and not to f(Y) in class T, even though t.f(Y) is a perfect match to the run-time types!

## Java Example



## Java Example - 2

```
X x = new X();
```

```
Y y = new Y();
```

```
X xy = (X) y;
```

x.f(x) -- invokes X:f(X)

x.f(y) -- invokes X:f(Y) since more specific than X:f(X)

x.f(xy) -- invokes X:f(X)

y.f(x) -- invokes Y:f(X), since more specific than X:f(X)

xy.f(x) -- invokes Y:f(X) which overrides X:f(X) for Y receivers.

