

## Computation to Core Mapping

- Lessons learned from a simple application


## A Simple Application

$>$ Matrix Multiplication
> Used as an example throughout the course
$>$ Goal for today:
$>$ Show the concept of "Computation-to-Core Mapping"
> Block schedule, Occupancy, and thread schedule

- Assumption
$>$ Deal with square matrix for simplicity
$>$ Leave memory issues later
> With global memory and local registers


## The algorithm and CPU code

## $\mathrm{P}=\mathrm{M}$ * N of size WIDTH $\times$ WIDTH

// Matrix multiplication on the (CPU) host in double precision void MatrixMulOnHost(float* M, float* N, float* P, int Width) \{
for (int i = 0; i < Width; ++i)
for (int j = 0; j < Width; ++j) \{
double sum $=0$; for (int k = 0; k < Width; ++k) \{ double $a=M[i *$ width $+k]$; double $\mathrm{b}=\mathrm{N}[\mathrm{k}$ * width +j$]$; sum += a * b; \}
$\mathrm{P}[\mathrm{i}$ * Width + j] = sum;
\}
\}

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$\mathrm{P}[\mathrm{i}$ * Width + j] = sum;
\}
\}


## First Mapping Scheme

> Thread mapping:
> Define the finest computational unit, and map it onto each thread
> Main criterion : None Dependency
> In our first scheme:
Unit: Calculation of one element of $P$
Block mapping:
> Simple: One block


## Step 1: Memory layout

| $M_{0,0}$ | $M_{1,0}$ | $M_{2,0}$ | $M_{3,0}$ |
| :--- | :--- | :--- | :--- |
| $M_{0,1}$ | $M_{1,1}$ | $M_{2,1}$ | $M_{3,1}$ |
| $M_{0,2}$ | $M_{1,2}$ | $M_{2,2}$ | $M_{3,2}$ |
| $M_{0,3}$ | $M_{1,3}$ | $M_{2,3}$ | $M_{3,3}$ |

$\mathbf{M}_{\text {(column\#, row\#) }}$
$M$
$\downarrow$

$$
\begin{array}{|l|l|l|l|l|l|l|l|l|l|l|l|l|l|}
\hline M_{0,0} & M_{1,0} & M_{2,0} & M_{3,0} & M_{0,1} & M_{1,1} & M_{2,1} & M_{3,1} & M_{0,2} & M_{1,2} & M_{2,2} & M_{3,2} & M_{0,3} & M_{1,3}
\end{array} M_{2,3} M_{3,3} .
$$

## Step 2: Input Matrix Data Transfer (Host Code)

void MatrixMulOnDevice(float* M , float* N , float* ${ }^{*}$, int Width)
\{
int size = Width * Width * sizeof(float);
float* Md, Nd, Pd;

1. // Allocate and Load $M, N$ to device memory cudaMalloc(\&Md, size);
cudaMemcpy(Md, M, size, cudaMemcpyHostToDevice);
cudaMalloc(\&Nd, size);
cudaMemcpy(Nd, N, size, cudaMemcpyHostToDevice);
// Allocate P on the device
cudaMalloc(\&Pd, size);

## Step 3: Output Matrix Data Transfer (Host Code)

2. // Kernel invocation code - to be shown later
3. // Read P from the device cudaMemcpy(P, Pd, size, cudaMemcpyDeviceToHost);
// Free device matrices
cudaFree(Md); cudaFree(Nd); cudaFree (Pd); \}

## Step 4: Kernel Function

// Matrix multiplication kernel - per thread code
_global__ void MatrixMulKernel(float* Md, float* Nd, float* Pd, int Width))
\{
// Pvalue is used to store the element of the matrix
// that is computed by the thread
float Pvalue $=0$;

## Step 4: Kernel Function (cont.)

$$
\begin{aligned}
& \text { for (int } k=0 ; k<\text { Width; ++k) \{ \{ } \\
& \text { float Melement }=\operatorname{Md}[\text { threadldx.y*Width+k]; } \\
& \text { float Nelement }=\operatorname{Nd}\left[k^{*}\right. \text { Width+threadldx.x]; } \\
& \text { Pvalue += Melement * Nelement; } \\
& \text { \} } \\
& \text { Pd[threadldx.y*Width+threadldx.x] = Pvalue; }
\end{aligned}
$$ \}



## Step 5: Kernel Invocation (Host Code)

// Setup the execution configuration dim3 $\operatorname{dim} \operatorname{Grid}(1,1)$; dim3 dimBlock(Width, Width);

## // Launch the device computation threads!

MatrixMulKernel<<<dimGrid, dimBlock>>>(Md, Nd, Pd, Width);

## Issues with the First Mapping Scheme

> One Block of threads compute matrix Pd
> Other Multi-processors are not used.


## Issues with the First Mapping Scheme

> Each thread
$>$ Loads a row of matrix Md
> Loads a column of matrix Nd
> Perform one multiply and addition for each pair of Md and Nd elements
$>$ Compute to off-chip memory access ratio close to 1:1 (not very high)


## Issues with the First Mapping Scheme

$>$ Size of matrix limited by the number of threads allowed in a thread block
> Maximum threads per block: 1024
> Can only do $32 \times 32$ matrix
$>$ You can put a loop around the kernel call for cases when Width > 32. But multiple kernel launch will cost you.


## Solution: the Second Mapping Scheme

> Thread mapping: the same with the first one $>$ Block mapping:
$>$ Create 2D thread blocks, each of which compute a (TILE_WIDTH) ${ }^{2}$ sub-matrix (tile) of the result matrix
$>$ Each has (TILE_WIDTH) ${ }^{2}$ threads
$>$ Generate a 2D Grid of (WIDTH/TILE_WIDTH) ${ }^{2}$ blocks


## About the Second Mapping

> More blocks (WIDTH/TLLE_WIDTH) ${ }^{2}$
> Support larger matrix
$>$ The maximum size of each dimension of a grid of thread blocks is 65535 .
$>$ Max Width $=65535 \times$ TILE_WIDTH
> Use more multi-processors



## Algorithm concept using tiles

$>$ Break-up Pd into tiles
> Each block calculates one tile
> Each thread calculates one element
> Block size equal tile size


## Example



## Block Computation



## Kernel Code using Tiles

```
__global__ void MatrixMulKernel(float* Md, float* Nd, float* Pd, int Width)
{
// Calculate the row index of the Pd element and M
int Row = blockIdx.y*TILE_WIDTH + threadIdx.y;
// Calculate the column idenx of Pd and N
int Col = blockIdx.x*TILE_WIDTH + threadIdx.x;
float Pvalue = 0;
// each thread computes one element of the block sub-matrix
for (int k = 0; k < Width; ++k)
    Pvalue += Md[Row*Width+k] * Nd[k*Width+Col];
Pd[Row*Width+Col] = Pvalue;
}
```


## Revised Kernel Invocation (Host Code)

// Setup the execution configuration dim3 dimGrid (Width/TILE_WIDTH, Width/TILE_WIDTH); dim3 dimBlock (TILE_WIDTH, TILE_WIDTH);
// Launch the device computation threads!
MatrixMulKernel<<<dimGrid, dimBlock>>>(Md, Nd, Pd, Width);

## Questions?

$>$ For Matrix Multiplication using multiple blocks, should I use 8X8, 16X16 or 32X32 blocks?
$>$ Why?

## Block Scheduling


$>$ Could be 256 (threads/block) * 3 blocks
$>$ Or 128 (threads/block) * 6 blocks, etc.
$>$ SM in GT200 can take up to 1024 threads

## Thread scheduling in Multiprocessing


> Each Block is executed as 32thread Warps
> If 3 blocks are assigned to an SM and each block has 256 threads, how many Warps are there in an SM?
$>$ Each Block is divided into 256/32 = 8 Warps

> There are 8 * $3=24$ Warps

## Occupancy of Multiprocessor

> How much a Multiprocessor is occupied: Occupancy = Actually warps / Totally allowed
$>$ GF 100 SM allows 48 warps

- GT200 SM allows 32 warps
$>$ G80 SM allow 24 warps
$>$ For example:
$>$ One block per SM, 32 threads per block $>(32 / 32) / 32=3.125 \%$ (Very bad)
$>4$ blocks per SM, 256 threads per block
$>(256 / 32) * 4 / 32=100 \%$ (Very good)
$>$ There are three factors:
$>$ Maximum number of warps
$>$ Maximum registers usage
$>$ Maximum share memory usage


## Answers to Our Questions

> For Matrix Multiplication using multiple blocks, should I use 8X8, 16X16 or 32X32 blocks?
> For G80 GPU:
> For 8X8, we have 64 threads per Block. Since each SM can take up to 768 threads, there are 12 Blocks. However, each SM can only take up to 8 Blocks, only 512 threads will go into each SM! (Occupancy = 66.6\%)
$>$ For 16X16, we have 256 threads per Block. Since each SM can take up to 768 threads, it can take up to 3 Blocks and achieve full capacity unless other resource considerations overrule. (Occupancy = 100\%)
$>$ For 32X32, we have 1024 threads per Block. Not even one can fit into an SM! (Can not support)

## Answers to Our Questions (Cont')

> For Matrix Multiplication using multiple blocks, should I use 8X8, 16X16 or 32X32 blocks?
$>$ For GT200 GPU:
> For 8X8, we have 64 threads per Block. Since each SM can take up to 1024 threads, there are 16 Blocks. However, each SM can only take up to 8 Blocks, only 512 threads will go into each SM! (Occupancy =50\%)
$>$ For 16X16, we have 256 threads per Block. Each SM takes 4 Blocks and achieve full capacity unless other resource considerations overrule. (Occupancy = 100\%)
$>$ For 32X32, we have 1024 threads per Block. Each SM takes 1 Block and achieve full capacity unless other resource considerations overrule. (Occupancy = 100\%)

## Computation-to-Core Mapping

$>$ Step 1:
$>$ Define your computational unit, map each unit to a thread
$>$ Avoid dependency
$>$ Increase compute to memory access ratio
$>$ Step 2:
> Group your threads into blocks
> Eliminate hardware limit
$>$ Take advantage of shared memory

