CS2104 Problem Solving in Computer Science

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Python and Algorithms







Day 1 Objectives

- Explain algorithms and efficiency
- Interpret flow charts
- Implement algorithms in python
- Experiment using LLMS to code and test





Algorithms

Algorithms are the threads that tie together most of the subfields of computer science.

Something magically beautiful happens when a sequence of commands and decisions is able to marshal a collection of data into organized patterns or to discover hidden structure.

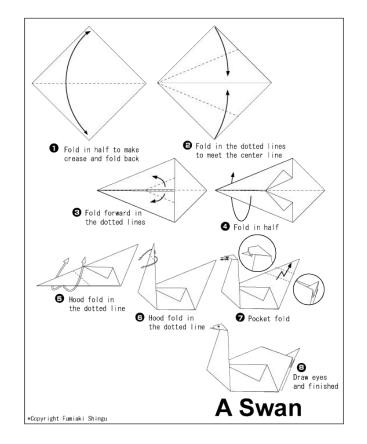
Donald Knuth







Example: Diagram for making an origami swan







Examples of Algorithms

- Recipe for Chicken and Macaroni Casserole
- Knitting pattern for a blanket
- Calculating a tip at a restaurant
- Finding the best deal on EXPO markers
- Finding the fastest route to Virginia Beach





Definition of Algorithm

effective method (or procedure)

- a procedure that reduces the solution of some class of problems to a series of rote steps which, if followed to the letter, and as far as may be necessary, is bound to:
 - always give some answer rather than ever give no answer;
 - always give the right answer and never give a wrong answer;
 - always be completed in a finite number of steps, rather than in an infinite number;
 - work for all instances of problems of the class.

Algorithm

 an effective method expressed as a finite list of well-defined instructions for calculating a function





Properties of an Algorithm

An algorithm must possess the following properties:

- **finiteness**: The algorithm must always terminate after a finite number of steps.
- definiteness: Each step must be precisely defined; the actions to be carried out must be rigorously and unambiguously specified for each case.
- **input:** An algorithm has zero or more inputs, taken from a specified set of objects.
- output: An algorithm has one or more outputs, which have a specified relation to the inputs.
- **effectiveness:** All operations to be performed must be sufficiently basic that they can be done exactly and in finite length.







Common **Elements** of Algorithms

Acquire data (input): Some means of reading values from an external source; most algorithms require data values to define the specific problem (e.g., coefficients of a polynomial)

Computation: Some means of performing arithmetic computations, comparisons, testing logical conditions, and so forth...

Selection: Some means of choosing among two or more possible courses of action, based upon initial data, user input and/or computed results

Iteration: Some means of repeatedly executing a collection of instructions, for a fixed number of times or until some logical condition holds

Report results (output): Some means of reporting computed results to the user, or requesting additional data from the user

Simple and list variables: Name and store data values





Modern CS Algorithms

- Matching users to servers, using Gayle-Shapely Algorithm for matching medical students to their residency placements
 - This is a beautiful algorithm for fair matching. Simple, elegant and effective. In its core form, it's also straightforward to implement. Has numerous applications. See: Stable marriage problem – Wikipedia
- Music Search using Fast Fourier Transforms (FFT)
 - Music recognition is done by converting it into frequency domain using FFT. FFT has implementations in number of languages. See this article for a great start: Shazam It! Music Recognition Algorithms, Fingerprinting, and Processing. © 2025 Ellis, Ramakrishnan & Nizamani — CC BY-NC-ND



Problems, Algorithms and Programs

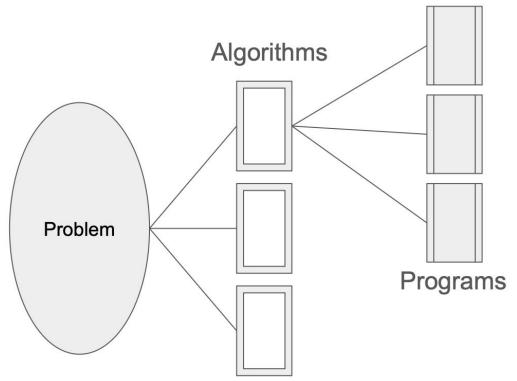
- For each problem or class of problems, there may be many different algorithms.
- For each algorithm, there may be many different implementations (programs).







Problems vs Algorithms vs Programs







Expressing Algorithms

An algorithm may be expressed in a number of ways:

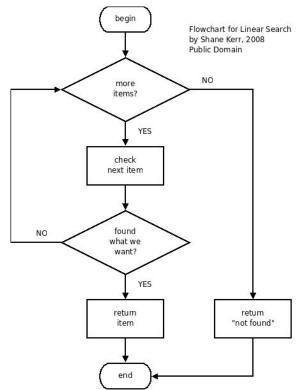
- natural language: usually verbose and ambiguous
- flow charts: avoid most (if not all) issues of ambiguity; difficult to modify w/o specialized tools; largely standardized
- pseudo-code: also avoids most issues of ambiguity; vaguely resembles common elements of programming languages; no particular agreement on syntax
- programming language: tend to require expressing low-level details that are not necessary for a high-level understanding







Flowchart Example of Linear Search







Testing Correctness

- How do we know whether an algorithm is actually correct?
- First, perform logical analysis
- Second, testing
- testing can never prove that the algorithm produces correct results in all cases.





Recap on Unit Test

- 1. Set up the initial conditions for your test (create any objects needed, place them in the correct state, put everything necessary where it needs to be, etc.).
- 2. Call the method you are testing.
- 3. Check that the behavior you expected has occurred. This could involve checking the return value of the method, or checking the state of the objects involved in the test. Be sure to check everything you expect to happen, not just the most obvious item.

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Polya's Problem Solving Steps

Understand your problem solving process then decide how and when an LLM can and cannot assist.

- Step 1: Understand the problem.
 - Have concrete examples (these become test cases)
- Step 2: Devise a plan (translate).
 - Break the problem down into manageable pieces/step
- Step 3: Carry out the plan (solve).
 - Iterate on 3 and 4 piece by piece
- Step 4: Look back (check and interpret).
 - Need to really understand the problem!



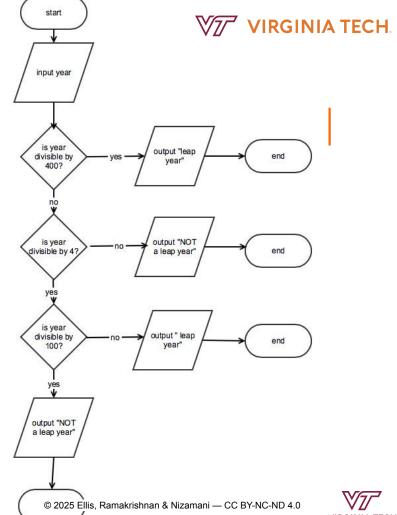


Class Activity -Turn Flowcharts into Python Code New Algorithm:

- 1. Turn Leap Year Flowchart into python code
- 2. Write test cases for each branch in the flowchart
- 3. Come up with alternate ways to code it



Leap Year Flowchart







Class Activity

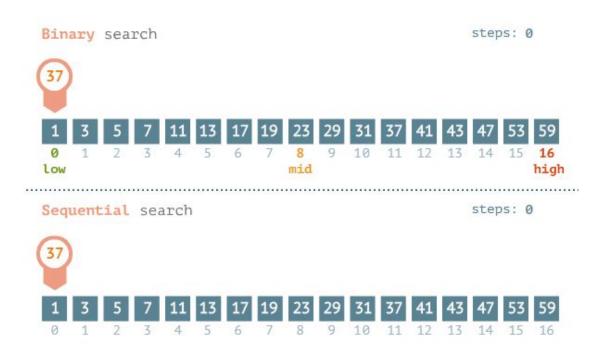
Do python first steps in class (use vscode in Rlogin)







Animation of Binary Search vs Sequential







Efficiency of Linear Search

- In many applications, it is easy to come up with a numeric value that specifies the problem size, which is generally denoted by the letter *N*.
- For most array applications, the problem size is simply the size of the array.
- In the worst case—which occurs when the value you're searching for comes at the end of the array or does not appear at all—linear search requires *N* steps.
- On average, it takes approximately half that time.
 - A linear search of 10 items takes an average of 5 looks
 - A linear search of 800 items takes an average of 400 looks





Efficiency of Binary Search

- The running time of binary search also depends on the number of elements, but in a profoundly different way.
- On each step in the process, the binary search algorithm rules out half of the remaining possibilities.
- In the worst case, the number of steps required is equal to the number of times you can divide the original size of the array in half until there is only one element remaining.





Efficiency of Binary Search (Cont.)

• In other words, what you need to find is the value of *k* that satisfies the following equation:

$$1=rac{N}{2 imes2 imes2 imes imes}$$
 $k ext{ times}$ $1=rac{N}{2^k}$ $2^k=N$ $k=\log_2 N$



Assessing Algorithmic Efficiency

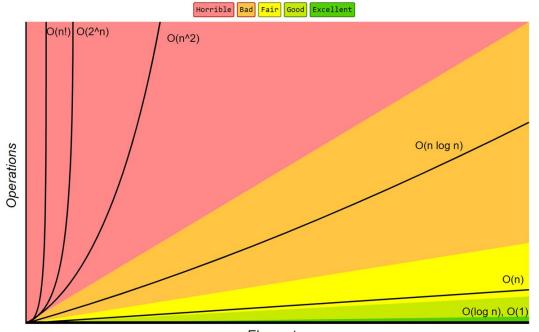
- The discussion of the efficiency of the various searching and sorting algorithms illustrates a fundamental computer science technique called **algorithmic analysis**.
- One of the most important problems in algorithmic analysis is deducing the **computational complexity** of an algorithm, which is the relationship between the size of the problem and the expected running time.





Examples of Algorithmic Efficiency

Big-O Complexity Chart







Big-O Notation

- The most common way to express computational complexity is to use **big-O notation**, which was introduced by the German mathematician Paul Bachmann in 1892.
- Big-O notation consists of the letter O followed by a formula that offers a qualitative assessment of running time as a function of the problem size, traditionally denoted as *N*.
- For example, the computational complexity of linear search is O(N) and the computational complexity of radix sort is $O(N \log N)$
- If you read these formulas aloud, you would pronounce them as "big-O of N" and "big-O of N log N" respectively.





Exercise: Computational Complexity

Assuming that none of the steps in the body of the following **for** loops depend on the problem size stored in the variable **n**, what is the computational complexity of each of the following examples:

```
a) for (int i = 0; i < n; i++) {
   for (int j = 0; j < i; j++) {
        ...loop body...
   }
}</pre>
```

```
b) for (int k = 1; k <= n; k *= 2) {
...loop body...
}
```

```
c) for (int i = 0; i < 100; i++) {
   for (int j = 0; j < i; j++) {
        ...loop body...
   }
}</pre>
```





Algorithm's Time Complexity

Big O Notation	Name	Example(s)
O(1)	Constant	# Odd or Even number, # Look-up table (on average)
O(log n)	Logarithmic	# Finding element on sorted array with binary search
O(n)	Linear	# Find max element in unsorted array, # Duplicate elements in array with Hash Map
O(n log n)	Linearithmic	# Sorting elements in array with merge sort
$O(n^2)$	Quadratic	# Duplicate elements in array **(naïve)**, # Sorting array with bubble sort
$O(n^3)$	Cubic	# 3 variables equation solver
O(2 ⁿ)	Exponential	# Find all subsets
O(n!)	Factorial	# Find all permutations of a given set/string



Binary Search

Given totals = {23, 44, 55, 146, 255, 300, 440, 575}

Trace binarySearch(440, totals)

List all the values of lh, rh and mid.

```
        0
        1
        2
        3
        4
        5
        6
        7

        23
        44
        55
        146
        255
        300
        440
        575
```

lh	rh	mid
0	7	3
4	7	5
6	7	6 🔽

def binarySearch (key,values):

```
lh = 0
rh = len(values) - 1
while (lh <= rh):
  mid = (lh + rh) // 2
  if (key == values[mid]):
     return mid
  if (key < values[mid]):
     rh = mid - 1
  else:
     lh = mid + 1
return - 1
```





LLMs to help (as allowed by assignments)

Sample Response: This unit test covers various scenarios, including cases where the key is present in the list, not present in the list, and a single-element list. You can run this test to verify that your binarySearch function works correctly. If you have any other questions or need further assistance, feel free to ask!



import

Sample Response Code:

import unittest

```
def binarySearch(key, values):
    Ih = 0
    rh = len(values) - 1

while lh <= rh:
    mid = (lh + rh) // 2
    if key == values[mid]:
        return mid
    if key < values[mid]:
        rh = mid - 1
    else:
        lh = mid + 1
    return -1....</pre>
```

```
#co-pilot 12/29/24
class TestBinarySearch(unittest.TestCase):
  def test binary search(self):
     # Test cases where the key is present in the list
     self.assertEqual(binarySearch(3, [1, 2, 3, 4, 5]), 2)
     self.assertEqual(binarySearch(1, [1, 2, 3, 4, 5]), 0)
     self.assertEqual(binarySearch(5, [1, 2, 3, 4, 5]), 4)
     # Test cases where the key is not present in the list
     self.assertEqual(binarySearch(0, [1, 2, 3, 4, 5]), -1)
     self.assertEqual(binarySearch(6, [1, 2, 3, 4, 5]), -1)
     # Test cases with an empty list
     self.assertEqual(binarySearch(3, []), -1)
     # Test cases with a single element list
     self.assertEqual(binarySearch(1, [1]), 0)
     self.assertEqual(binarySearch(0, [1]), -1)
if name == ' main ':
    unittest.main(argv=["], exit=False)
```

Unit test class





Find the Longest Run

- Given a list of values, finds the length of the longest sequence of values that are in strictly increasing order.
- Concrete example?
 - longest_run([5,4,8,1,2,3,1]) should output 3 as the longest run
 - What variables would you use?
 - What procedures would you take?
 - How do you test the correctness of your algorithm?





Find the Longest Run (Cont.)

```
# Given a list of values, finds the length of the longest sequence
# of values that are in strictly increasing order.
def longest run(arr):
  size = len(arr)
  if size \leq 0:
                                # if list is empty, no runs...
     print("The list is empty")
  else:
     current position = 0
                               # start with first element in list
     max run_length = 1
                               # it forms a run of length 1
     this run length = 1
```





Find the Longest Run (Continue)

```
# QUESTION TO PONDER: is this Algorithm Correct?
     while current position < size - 1:
       if (arr[current position] < arr[current position + 1]):
          this run length = this run length + 1
       else:
          if (this run length > max run length):
            max run length = this run length
                                                                       [0, 1, 2, 3]
          this run length = 1
                                                        [1, 3, 0, 5, -1, 1, 2, 3]
       current position = current position + 1
     print("The max run length is: ", max run length)
```



Unit Testing...

Demo Class Activity: Unit Testing,

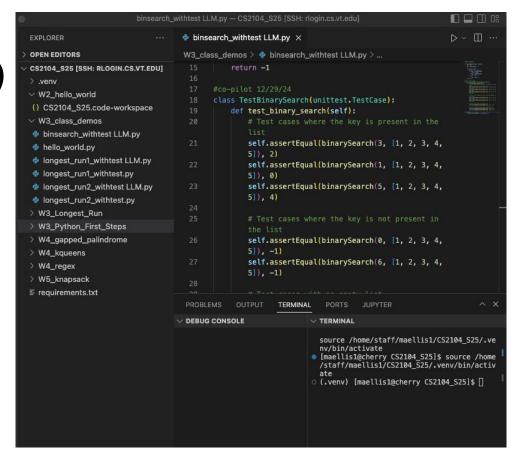
- students follow along
- create W3 demo folder in vscode
- click on the download arrow in canvas week 5 to get the py version of each file
 - binsearch_withtest_LLM.py
 - longest run1 withtest.py
 - longest run2 withtest.py
 - longest run1 withtest LLM.py
 - longest run2 withtest LLM.py
- Drag files into the directory (they get uploaded to rlogin)



Unit Testing (Cont.)

Demo Class Activity: Unit Testing,

binsearch withtest LLM.py should pass



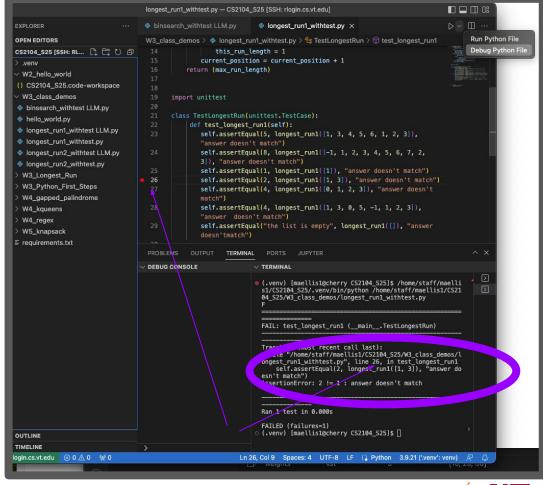
Unit Testing cont.

Demo Class Activity: Unit Testing

- Longest run1 withtest.py
 - should not pass
 - (investigate with debugger)
 - Set breakpoint in the gutter of failing test
 - Use buttons to trace code



- Watch values in variables
- Fixed version (can look side-by-side)
 Longest run2 withtest.py should pass









Unit Testing (Continue)

Demo Class Activity: Unit Testing (same code, different tests)

- longest run1 withtest LLM.py
 - Should not pass
 - Investigate with debugger
 - Use buttons to trace code(hover to see functionality)



- Watch values in variables
- longest run2 withtest LLM.py
 - Should pass
 - Compare side-by-side

```
♦ longest_run1_withtest LLM.py × ♦ longest !! ID ?

 RUN AND DE... ▷ No Conf ∨ 🚳 ···
                                     W3_class_demos > ♦ longest_run1_withtest LLM.py > ♦ longest_run1
VARIABLES
                                            def longest run1(arr):
V Locals
                                                if len(arr)==0:
                                                    return ("the list is empty")
   current_position: 5
   max_run_length: 3
                                                    current position = 0
   this_run_length: 4
                                                    max run length = 1
                                                    this run length = 1
  Globals
                                                while current_position < len(arr)-1:
                                                    if ( arr[current_position] < arr[current_position + 1]
                                                        this_run_length = this_run_length + 1
                                                        if (this run length > max run length):
                                                            max run length = this run length
                                                        this run length = 1
                                   D 15
                                                    current_position = current_position + 1
                                                return (max_run_length)
WATCH
                                            import unittest
```





Maximum Subsequence Sum Problem

• Given (possibly negative) integers A₁, A₂, ..., A_N, find the maximum value of:

$$\sum_{k=1}^{j} \left(A_k \right)$$

- The maximum subsequence is defined as zero if all the integers are negative, (i..e., the subsequence of the empty set is zero).
- Consider the sequence: 4 -3 5 -2 -1 2 6 -2 © 2025 Ellis, Ramakrishnan & Nizamani CC BY-NC-ND





Maximum subsequence sum solution 1

```
# finds the max subsequence sum in the list values
def max_subsequence_sum_1(values):
  size = len(values)
  max sum = 0
  i = 0
  while i < size:
    i = i
    while j < size:
       this sum = 0
       k = i
                                                        If size = 1000, how many
       while k \le j:
                                                        times are the statements in
         this_sum = this_sum + values[k]
                                                        the innermost loop
         k = k + 1
                                                        executed?
       if (this sum > max sum ):
         max sum = this sum
      j = j + 1
                                                      Can you/we do better?
    i = i + 1
  return(max_sum)
```



Maximum subsequence sum solution 2

 Reuse the subsum, so just adds the newest value in the subsequence to what was previously calculated

```
[4]
[4 -3]
[4 -3 5]
[4 -3 5]
[4 -3 5 -2]
[4 -3 5 -2 -1]
[4 -3 5 -2 -1 2]
[4 -3 5 -2 -1 2 6]
[4 -3 5 -2 -1 2 6 -2]
[5 5 -2 -1 2 6 -2]
```



Maximum subsequence sum solution 2 (Continue)

```
# finds the max subsequence sum in the list values
def max subsequence sum 2(values):
  size = len(values)
  max sum = 0
  i = 0
  while i < size:
    j = i
    this sum = 0
    while j < size:
      this_sum = this_sum + values[j]
      if (this_sum > max_sum ):
         max_sum = this_sum
      j = j + 1
  return(max sum)
```

If size = 1000, how many times are the statements in the innermost loop executed?

Can you/we do even better?





Maximum subsequence sum solution 3

```
# finds the max subsequence sum in the list values
def max_subsequence_sum_3(values):
  size = len(values)
  \max sum = 0
  this sum = 0
  i = 0
  while (i \le size - 1):
     this sum = this sum + values[j]
     if (this sum > max sum):
       \max \text{ sum} = \text{this sum}
     elif (this sum < 0):
       this sum = 0
    j = j + 1
  return max sum
```

Try it:

4 -3 5 -2 -1 2 6 -2

One must observe that if any list[i] is negative then it cannot be the beginning of the optimal sequence since any sequence starting with it would be improved by omitting it and starting with list[i+1].

Also by the same logic any negative subsequence cannot be the start of the optimal subsequence.

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Maximum subsequence sum solution LLM

```
def max_subsequence_sum(arr):
  max sum = 0
  current sum = 0
  for num in arr:
    current sum += num
    if current sum > max sum:
       max sum = current sum
    if current sum < 0:
                                         {result}")
       current sum = 0
  return max sum
```

```
# Given sequence
sequence = [4, -3, 5, -2, -1, 2, 6, -2]

# Calculate the maximum subsequence sum
result = max_subsequence_sum(sequence)

print(f"The maximum subsequence sum is:
{result}")
```



Research on LLMS in CS Education

- Students using LLMS had reduced ability to write code from scratch but similar ability to trace and read code
 - source: <u>CS1-LLM</u>: <u>Integrating LLMs into CS1 Instruction</u>
- Students who used AI less or later in problem solving process had higher grades in course (don't know cause-effect)
 - source: Interactions with How Novices Use LLMs to Solve Programming Problems
- Students with strong programming foundations benefited more from LLMs, while those without a solid base may be hindered by them
 - Sources: Insights from Social Shaping Theory: The Appropriation of Large Language Models in an Undergraduate Programming Course and The Widening Gap: The Benefits and Harms of Generative AI for Novice Programmers





Leetcode Assignment

- 1. Looks over assignment together
- 2. Open up Leetcode together(set to Python)
- 3. VT approved LLM (bing.com/chat)
- 4. Considerations
 - Can also experiment in vscode
 - Can also create your own test cases
 - Keep trying to get more efficient solutions

